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JUDGE CEDARBAUM

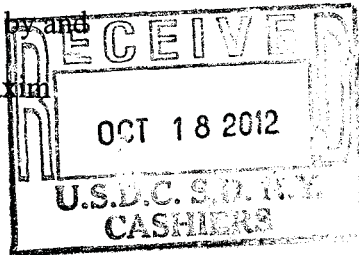
UNITED STATES DISTRICT COURT  
SOUTHERN DISTRICT OF NEW YORK

-----X  
Deutsche Bank AG New York Branch,  
Plaintiff,  
vs.  
Maxim Integrated Products, Inc.,  
Defendant.  
-----X

: No. \_\_\_\_\_  
: ECF CASE  
: **COMPLAINT FOR DECLARATORY**  
: **JUDGMENT**  
: **JURY TRIAL DEMANDED**

**COMPLAINT FOR DECLARATORY JUDGMENT**

1. Plaintiff Deutsche Bank AG New York Branch ("Deutsche Bank"), by and through its undersigned counsel, hereby files this Complaint against Defendant Maxim Integrated Products, Inc. ("Maxim") and avers as follows:



**NATURE OF THE ACTION**

2. This is an action based on the Patent Laws of the United States, 35 U.S.C. § 100 *et seq.*, seeking a declaratory judgment of non-infringement and/or invalidity with respect to United States Patent Nos. 5,940,510 (the "'510 patent'"), 5,949,880 (the "'880 patent'"), 6,105,013 (the "'013 patent'"), and 6,237,095 (the "'095 patent'") (collectively, the "'510, '880, '013, and '095 patents" or "the patents-in-suit").

**THE PARTIES**

3. Plaintiff Deutsche Bank AG New York Branch is a branch of Deutsche Bank AG, a German corporation. Deutsche Bank AG New York Branch is licensed by the State of New York Department of Financial Services and is located at 60 Wall Street, New York, New York 10005, within this District.

4. Defendant Maxim Integrated Products ("Maxim") is, on information and belief, a Delaware corporation with a place of business at 160 Rio Robles, San Jose, California 95134.

## **JURISDICTION AND VENUE**

5. This declaratory judgment claim arises under the United States Patent Laws, 35 U.S.C. § 100 et seq., and the Declaratory Judgment Act, 28 U.S.C. §§ 2201 and 2202. Subject matter jurisdiction is proper under 28 U.S.C. §§ 1331 and 1338(a). This matter presents an actual case or controversy and serves the purpose of resolving the legal rights of the parties.

6. This Court has personal jurisdiction over Maxim at least because Maxim has maintained continuous and systematic contacts with the State of New York, and Maxim has purposefully availed itself of the benefits and protections of the laws of the State of New York.

7. Maxim, directly and/or through its distribution and sales networks, offers for sale, sells, and/or distributes products within the State of New York.

8. Maxim maintains sales offices within the State of New York. *See* Exhibit A, a printout from Maxim's website <http://www.maximintegrated.com/sales/offices/worldwide.mvp>, showing its worldwide sales offices, including one in New York.

9. Maxim engages franchised distributors within the State of New York. *See* Exhibit B, a printout from Maxim's website <http://www.maximintegrated.com/sales/offices/distributor/franchise.mvp>, showing four franchised distributors in New York.

10. In addition, Maxim operates an interactive website through which persons in New York can and do order products from Maxim, which are shipped to New York. *See* <http://www.maximintegrated.com/sales/>.

11. Venue is proper in this judicial district under at least 28 U.S.C. §§ 1391(b) and (c).

## **FACTUAL BACKGROUND**

12. On or about December 22, 2011, counsel for Maxim sent a letter that included the following statements:

Maxim has analyzed Deutsche Bank's products and services within the mobile transaction space. It is our belief that Deutsche Bank is infringing a number of the patents within the Maxim Mobile Transaction Patent Portfolio. In particular, we believe that the Deutsche Bank mobile platforms, including its multiple software applications for the iPhone and other mobile devices, infringe certain claims within the portfolio via direct infringement, joint infringement, contributory infringement and/or inducement. A technical analysis document of one of the Deutsche Bank mobile applications and exemplary claim charts are provided herewith.

13. The letter that Maxim's counsel sent on or about December 22, 2011, identified the following patents as being part of the "Maxim Mobile Transaction Patent Portfolio":

## **APPENDIX I**

### **Maxim Mobile Transaction Patent Portfolio**

#### **General Secure Transaction Patents**

1. U.S. Patent No. 5,940,510: "Transfer of Valuable Information Between a Secure Module and Another Module"
2. U.S. Patent No. 5,949,880: "Transfer of Valuable Information Between a Secure Module and Another Module"
3. U.S. Patent No. 6,105,013: "Method, Apparatus, System and Firmware for Secure Transactions"
4. U.S. Patent No. 6,237,095: "Apparatus for Transfer of Secure Information Between a Data Carrying Module and an Electronic Device"

14. True and correct copies of the '510, '880, '013, and '095 patents are attached hereto as Exhibits C, D, E, and F.

15. Upon information and belief, the '510, '880, '013, and '095 patents are assigned to Maxim.

16. The "exemplary claim charts" referenced in the letter that Maxim's counsel sent on or about December 22, 2011, asserted infringement of at least:

- Claim 9 of the '013 patent
- Claim 1 of the '095 patent
- Claim 1 of the '880 patent
- Claim 1 of the '510 patent

17. The “technical analysis document” referenced in the letter that Maxim’s counsel sent on or about December 22, 2011, included the following statements:

Maxim Integrated Products has analyzed the DEUTSCHE mobile banking software application to objectively show infringement of certain patents within its Secure Mobile Transaction Patent Portfolio. This analysis included downloading the mobile software, analyzing packets exchanged between the mobile handset and server, and de-encrypting network traffic transmitted from the server that establishes agreements related to user authentication, encryption, randomization, time-stamping, etc. Based on this analysis, we are able to objectively show that the DEUTSCHE bank infringes certain claims within Maxim patents. ***This document is referenced by the claim charts, provided herewith, showing the DEUTSCHE mobile banking application, server structures and processes, and overall system architecture infringe a diverse set of claims within the Maxim patent portfolio.***

18. The letter that Maxim’s counsel sent on or about December 22, 2011, also stated, “If I do not hear back from Deutsche Bank within the one month period, Maxim will assume that Deutsche Bank does not want to obtain a license in a non-litigious manner and will act accordingly.”

19. Maxim’s letter sent on or about December 22, 2011, together with the “technical analysis document” and “exemplary claim charts” alleging patent infringement, constitute a clear and unmistakable threat of litigation against Deutsche Bank’s mobile banking software.

20. Within the past year, Maxim has filed numerous lawsuits against different entities alleging infringement of the ’510, ’880, ’013, and ’095 patents.

21. On or about January 6, 2012, Maxim filed separate patent infringement lawsuits against (1) Capital One Financial Corporation, (2) Expedia, Inc., and (3) Starbucks Corporation, respectively.

22. On or about January 9, 2012, Maxim filed a patent infringement lawsuit against (4) Bank of the West.

23. On or about January 11, 2012, Maxim filed a patent infringement lawsuit against (5) First United Bank & Trust Co.

24. On or about February 22, 2012, Maxim filed separate patent infringement lawsuits against (6) QVC, Inc., (7) Union Bank, N.A., and (8) Southwest Airlines Co., respectively.

25. On or about February 23, 2012, Maxim filed separate patent infringement lawsuits against (9) Groupon, Inc. and (10) Comerica Inc., respectively.

26. On or about October 1, 2012, Maxim filed separate patent infringement lawsuits against (11) Bank of America Corporation, (12) CitiGroup, Inc., (13) JPMorgan Chase & Co., (14) Target Corporation, (15) U.S. Bancorp, (16) Wal-Mart Stores, Inc., and (17) Wells Fargo & Co.

27. In summary, Maxim has filed at least 17 patent infringement lawsuits within the past year asserting infringement of the '510, '880, '013, and '095 patents, and at least 10 of those lawsuits have been against banks and/or financial institutions.

28. Based on at least the foregoing facts, there is a substantial controversy between Maxim and Deutsche Bank, of sufficient immediacy and reality, to warrant the issuance of a declaratory judgment with respect to non-infringement and/or invalidity of at least claim 9 of the '013 patent, claim 1 of the '095 patent, claim 1 of the '880 patent, and claim 1 of the '510 patent.

**COUNT I: DECLARATORY JUDGMENT OF NON-INFRINGEMENT**  
**OF THE '013 PATENT**

29. Deutsche Bank incorporates the preceding paragraphs as if fully set forth herein.

30. Deutsche Bank has not infringed, and does not infringe, at least claim 9 of the '013 patent.

31. Therefore, Deutsche Bank is entitled to a declaration of non-infringement with respect to one or more claims of the '013 patent.

**COUNT II: DECLARATORY JUDGMENT OF INVALIDITY**  
**OF THE '013 PATENT**

32. Deutsche Bank incorporates the preceding paragraphs as if fully set forth herein.

33. In the alternative and/or in addition, at least claim 9 of the '013 patent is invalid.

34. Therefore, in the alternative and/or in addition, Deutsche Bank is entitled to a declaration of invalidity with respect to one or more claims of the '013 patent.

**COUNT III: DECLARATORY JUDGMENT OF NON-INFRINGEMENT**  
**OF THE '095 PATENT**

35. Deutsche Bank incorporates the preceding paragraphs as if fully set forth herein.

36. Deutsche Bank has not infringed, and does not infringe, at least claim 1 of the '095 patent.

37. Therefore, Deutsche Bank is entitled to a declaration of non-infringement with respect to one or more claims of the '095 patent.

**COUNT IV: DECLARATORY JUDGMENT OF INVALIDITY**  
**OF THE '095 PATENT**

38. Deutsche Bank incorporates the preceding paragraphs as if fully set forth herein.

39. In the alternative and/or in addition, at least claim 1 of the '095 patent is invalid.

40. Therefore, in the alternative and/or in addition, Deutsche Bank is entitled to a declaration of invalidity with respect to one or more claims of the '095 patent.

**COUNT V: DECLARATORY JUDGMENT OF NON-INFRINGEMENT**  
**OF THE '880 PATENT**

41. Deutsche Bank incorporates the preceding paragraphs as if fully set forth herein.

42. Deutsche Bank has not infringed, and does not infringe, at least claim 1 of the '880 patent.

43. Therefore, Deutsche Bank is entitled to a declaration of non-infringement with respect to one or more claims of the '880 patent.

**COUNT VI: DECLARATORY JUDGMENT OF INVALIDITY**  
**OF THE '880 PATENT**

44. Deutsche Bank incorporates the preceding paragraphs as if fully set forth herein.

45. In the alternative and/or in addition, at least claim 1 of the '880 patent is invalid.

46. Therefore, in the alternative and/or in addition, Deutsche Bank is entitled to a declaration of invalidity with respect to one or more claims of the '880 patent.

**COUNT VII: DECLARATORY JUDGMENT OF NON-INFRINGEMENT**  
**OF THE '510 PATENT**

47. Deutsche Bank incorporates the preceding paragraphs as if fully set forth herein.

48. Deutsche Bank has not infringed, and does not infringe, at least claim 1 of the '510 patent.

49. Therefore, Deutsche Bank is entitled to a declaration of non-infringement with respect to one or more claims of the '510 patent.

**COUNT VIII: DECLARATORY JUDGMENT OF INVALIDITY**  
**OF THE '510 PATENT**

50. Deutsche Bank incorporates the preceding paragraphs as if fully set forth herein.

51. In the alternative and/or in addition, at least claim 1 of the '510 patent is invalid.

52. Therefore, in the alternative and/or in addition, Deutsche Bank is entitled to a declaration of invalidity with respect to one or more claims of the '510 patent.

**PRAYER FOR RELIEF**

WHEREFORE, Deutsche Bank respectfully requests that this Court enter judgment in its favor as follows:

- a) declare that Deutsche Bank has not infringed, and does not infringe, one or more claims of the '510, '880, '013, and '095 patents;
- b) declare that one or more claims of the '510, '880, '013, and '095 patents are invalid;
- c) declare this to be an exceptional case and award Deutsche Bank its costs, expenses, and disbursements in this action, including reasonable attorney fees, pursuant to 35 U.S.C. § 285; and
- d) award Deutsche Bank any further and additional relief that this Court deems just and proper.

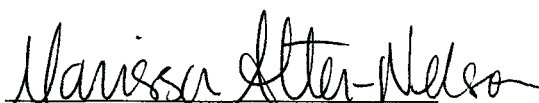
**DEMAND FOR JURY TRIAL**

53. Deutsche Bank hereby requests a trial by jury on all issues so triable.

Dated: October 18, 2012

Respectfully submitted,

SIDLEY AUSTIN LLP

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*Attorneys for Plaintiff  
Deutsche Bank AG New York Branch*



# **EXHIBIT A**



Maxim > Sales > Maxim Sales Offices

## Maxim Sales Offices

### Sales Inquiries

Maxim Direct: 888-MAXIM-IC (888-629-4642)

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- **Maxim Direct**
- **Worldwide Franchised Distributors**

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Austria	Maxim Direct Sales	0800 296 866 (toll-free) +353 1223 5501	+353 1223 5600
Belgium	Maxim Direct	0800 79145	+353 1223 5600
Brazil	DHW (Sales Representative)	+55 48 32255052	
Canada	Maxim Integrated	800-416-7237	
Canada	Sygnum Technologies (Sales Representative)	1-855-779-7466	
Chile	DHW (Sales Representative)	+55 48 32255052	
China	Maxim China	(86) 10 62115199	(86) 10 62115299
China - Shanghai	Dallas Semiconductor Asia Ltd.	(86) 21 68765001	(86) 21 58200188
Denmark	Maxim Direct	(45) 89292000 80 881 057 (toll-free)	(45) 89292020
France	Maxim Direct	0800 91 82 10 (toll-free) +33 (0)1 39 30 19 00	+33 (0) 1 39 30 19 10
Germany	Maxim Direct	0800 1824 943 (toll-free) +353 1223 5501	+353 1223 5600
India	Maxim India	+91-80-3028 5800	+91-80-3028 5812
Ireland	Maxim Ireland	+353 (0) 1 223 5500	+353 (0) 1 223 5600
Italy	Maxim Italy	(39) 039 64112	(39) 039 605 8566
Japan	Maxim Japan Co., Ltd.	+81 3 6893 6600	+81 3 6893 6935
Japan	Maxim Japan Osaka Office	+81 6 4704 3121	+81 6 4704 3128
Korea	Maxim Korea	+82 2 551 1010	+82 2 551 1218 +82 2 551 1219
Malaysia	Maxim Integrated	(65) 6841 7117	(65) 6841 7227
Mexico	Everest Sales and Solutions (Sales Representative)	+52 (33) 3123-0848	+52 (33) 3123-0847
Netherlands	Maxim Direct	0800-022 07 34 (toll-free)	+353 1223 5600
Norway	Maxim Direct	800 10247 (toll-free)	+353 1223 5600
Peru	DHW (Sales Representative)	+55 48 32255052	
Singapore	Maxim Integrated	(65) 6841 7117	(65) 6841 7227
Spain	Maxim Direct	900 972 932 (toll-free)	+353 1223 5600
Sweden	Maxim	020 791 801 (toll-free) 08 444 5430 +353 1223 5504	+353 1223 5600
Switzerland	Maxim Direct	0800 838 703 (toll-free) +41-44-828 80 80	+41-44-828 80 99
Taiwan	Maxim Taiwan	(886) 2 2547 3888	(886) 2 2547 3815
Thailand	Maxim Integrated	(65) 6841 7117	(65) 6841 7227

United Kingdom	Maxim Integrated (U.K.), Ltd.	+44 (0) 118 900 6300	+44 (0) 118 900 6400
US	Maxim North America Sales	888-MAXIM-IC (888-629-4642) or 972-371-6824	800-992-1884
US - Alabama	Maxim U.S. Sales	800-570-9738	
US - Arizona	Maxim U.S. Sales	800-570-9732	
US - California	Maxim U.S. Sales	408-737-0214 800-570-9732 (toll-free)	
US - Colorado	Maxim U.S. Sales	408-737-0214	
US - Colorado	Innovatech Associates Rocky Mountains (Sales Representative)	303-277-0456	303-974-1038
US - Connecticut	Maxim U.S. Sales	800-416-7237	
US - Delaware	Maxim U.S. Sales	800-416-7237	
US - Florida	Maxim U.S. Sales	800-570-9738	
US - Florida	Semtronic Associates (Sales Representative)	407-831-8233	407-804-0402
US - Georgia	Maxim U.S. Sales	800-570-9738	
US - Idaho	Maxim U.S. Sales	408-737-0214	
US - Idaho	Innovatech Associates Rocky Mountains (Sales Representative)	303-277-0456	303-974-1038
US - Illinois	Maxim U.S. Sales	800-570-9735	
US - Indiana	Maxim U.S. Sales	800-570-9735	
US - Indiana	Millennium Technical Sales (Sales Representative)	440-349-6600	
US - Iowa	Maxim U.S. Sales	800-570-9735	
US - Kansas	Maxim U.S. Sales	800-570-9735	
US - Kentucky	Millennium Technical Sales (Sales Representative)	440-349-6600	
US - Kentucky	Maxim U.S. Sales	800-570-9735	
US - Louisiana	Maxim U.S. Sales	800-631-1070	
US - Maine	Maxim U.S. Sales	800-416-7237	
US - Maryland	Maxim U.S. Sales	800-416-7237	
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US - Michigan	Maxim U.S. Sales	800-570-9735	
US - Michigan	Millennium Technical Sales (Sales Representative)	440-349-6600	
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US - Missouri	Maxim U.S. Sales	800-570-9735	
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US - Montana	Innovatech Associates Rocky Mountains (Sales Representative)	303-277-0456	303-974-1038
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US - New York	Maxim U.S. Sales	800-416-7237	
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US - Ohio	Millennium Technical Sales (Sales Representative)	440-349-6600	
US - Ohio	Maxim U.S. Sales	800-570-9735	
US - Oklahoma	Maxim U.S. Sales	800-631-1070	
US - Oregon	Maxim U.S. Sales	408-737-0214	
US - Pennsylvania	Maxim U.S. Sales	800-416-7237	
US - Pennsylvania (Western)	Millennium Technical Sales (Sales Representative)	440-349-6600	
US - Rhode Island	Maxim U.S. Sales	800-416-7237	
US - South Carolina	Maxim U.S. Sales	800-570-9738	
US - Texas	Maxim U.S. Sales	800-631-1070	512-249-2942
US - Utah	Maxim U.S. Sales	408-737-0214	

US - Utah	Innovatech Associates Rocky Mountains (Sales Representative)	303-277-0456	303-974-1038
US - Vermont	Maxim U.S. Sales	800-416-7237	
US - Virginia	Maxim U.S. Sales	800-416-7237	
US - Washington	Maxim U.S. Sales	408-737-0214	
US - Wisconsin	Maxim U.S. Sales	800-570-9735	
US - Wyoming	Innovatech Associates Rocky Mountains (Sales Representative)	303-277-0456	303-974-1038
Venezuela	DHW (Sales Representative)	+55 48 32255052	

# **EXHIBIT B**

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# Worldwide Franchised Distributors

Region: North America Country: United States State/Province: New York

	Name	State/Province	City	Phone	Fax
	Avnet	New York	Hauppauge	631-582-7700	
	Avnet	New York	Rochester	585-350-2800	
	Digi-Key	New York		800-344-4539	218-681-3380
	Mouser	New York		(800) 346-6873	(817) 804-3899

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# **EXHIBIT C**



US005940510A

# United States Patent [19]

[11] Patent Number: **5,940,510**

Curry et al.

[45] Date of Patent: **\*Aug. 17, 1999**

[54] **TRANSFER OF VALUABLE INFORMATION BETWEEN A SECURE MODULE AND ANOTHER MODULE**

[75] Inventors: **Stephen M. Curry**, Dallas; **Donald W. Loomis**, Coppel; **Michael L. Bolan**, Dallas, all of Tex.

[73] Assignee: **Dallas Semiconductor Corporation**, Dallas, Tex.

[\*] Notice: This patent issued on a continued prosecution application filed under 37 CFR 1.53(d), and is subject to the twenty year patent term provisions of 35 U.S.C. 154(a)(2).

[21] Appl. No.: **08/594,975**

[22] Filed: **Jan. 31, 1996**

[51] Int. Cl.<sup>6</sup> ..... **H04L 9/00**

[52] U.S. Cl. .... **380/25; 380/49**

[58] Field of Search ..... 380/49, 24, 23, 380/25

## [56] References Cited

### U.S. PATENT DOCUMENTS

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5,150,407	9/1992	Chan	380/4
5,189,700	2/1993	Blandsord	380/23
5,241,599	8/1993	Bellovin et al.	380/21
5,539,825	7/1996	Akiyama et al.	380/24
5,539,828	7/1996	Davis	380/50
5,546,463	8/1996	Caputo et al.	380/25
5,577,121	11/1996	Davis et al.	380/24
5,621,796	4/1997	Davis et al.	380/24
5,787,174	7/1998	Tuttle	380/23

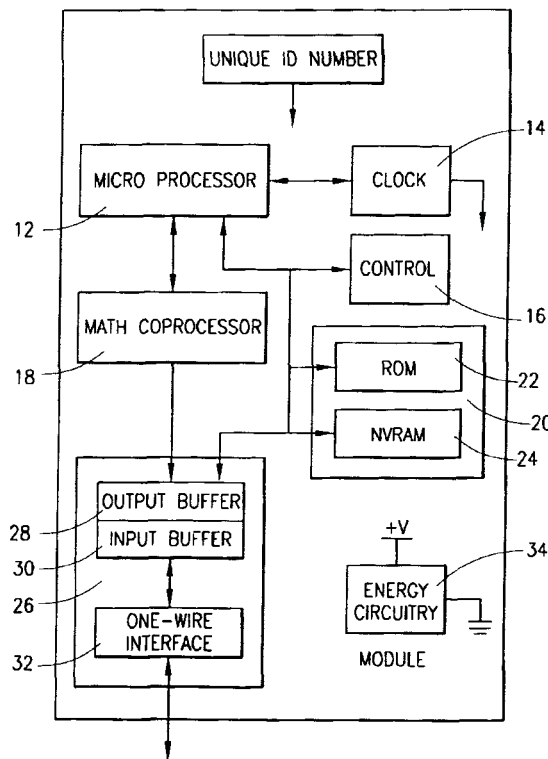
*Primary Examiner*—Salvatore Cangialosi  
*Attorney, Agent, or Firm*—Jenkins & Gilchrist

## [57] ABSTRACT

The present invention rotates to system, apparatus and method for communicating valuable data from a portable module to another module via an electronic device. More specifically, the disclosed system, apparatus and method are useful for enabling a user to fill a portable module with a cash equivalent and to spend the cash equivalent at a variety of locations. The disclosed system incorporates an encryption/decryption method.

**6 Claims, 8 Drawing Sheets**

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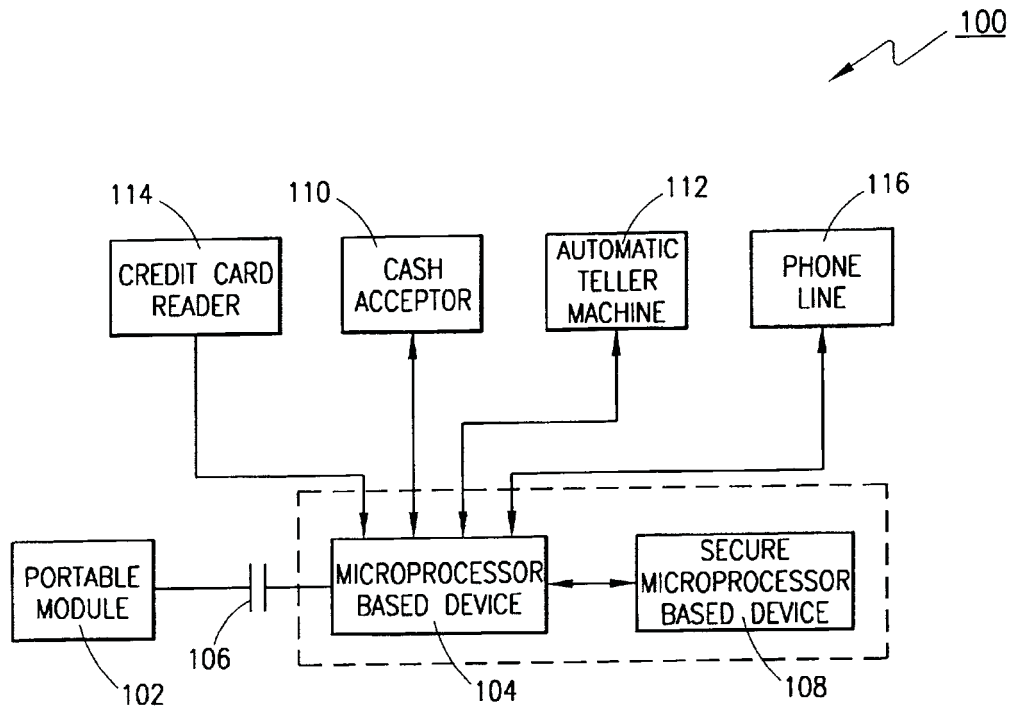


FIG. 1

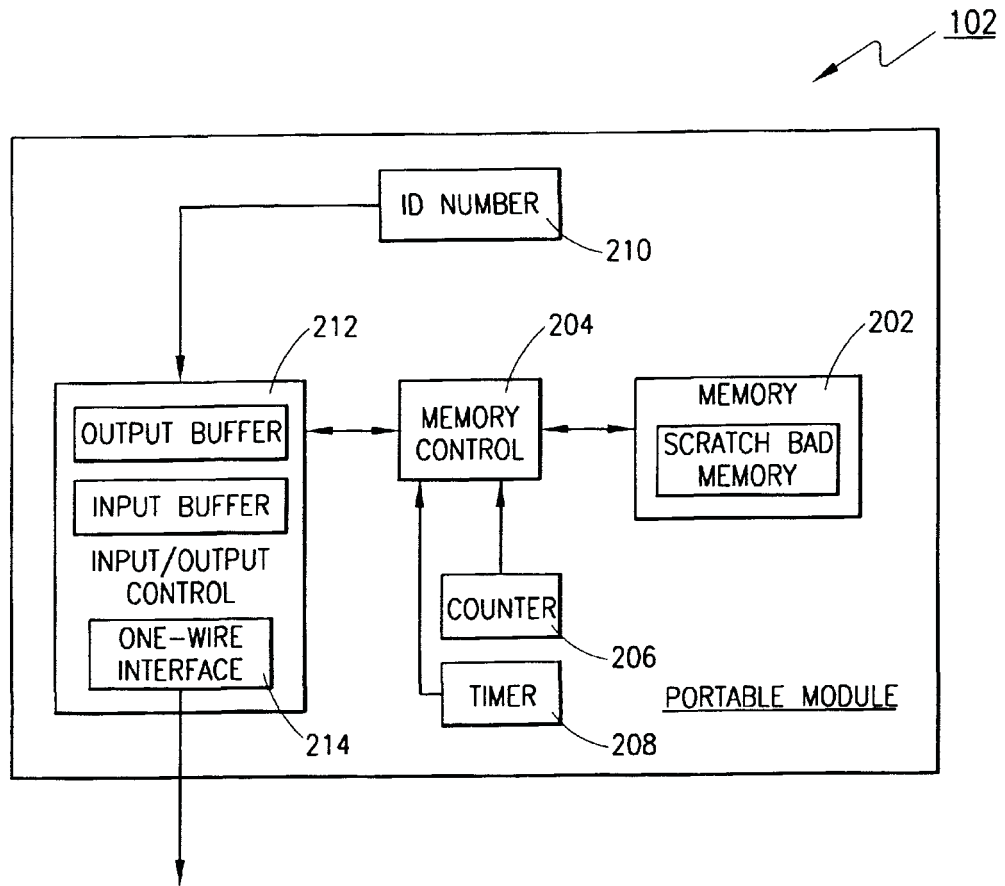


FIG. 2

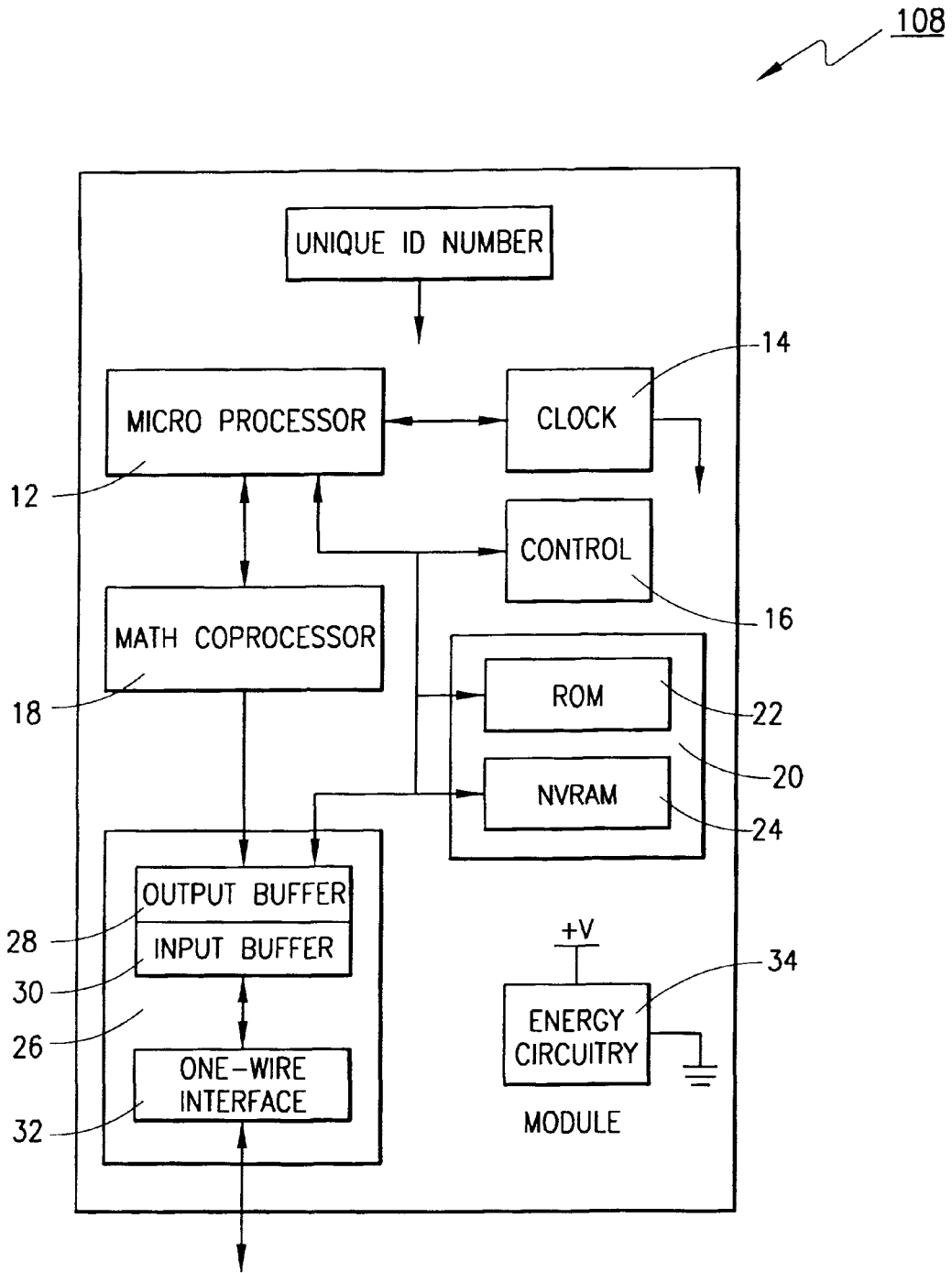


FIG. 3

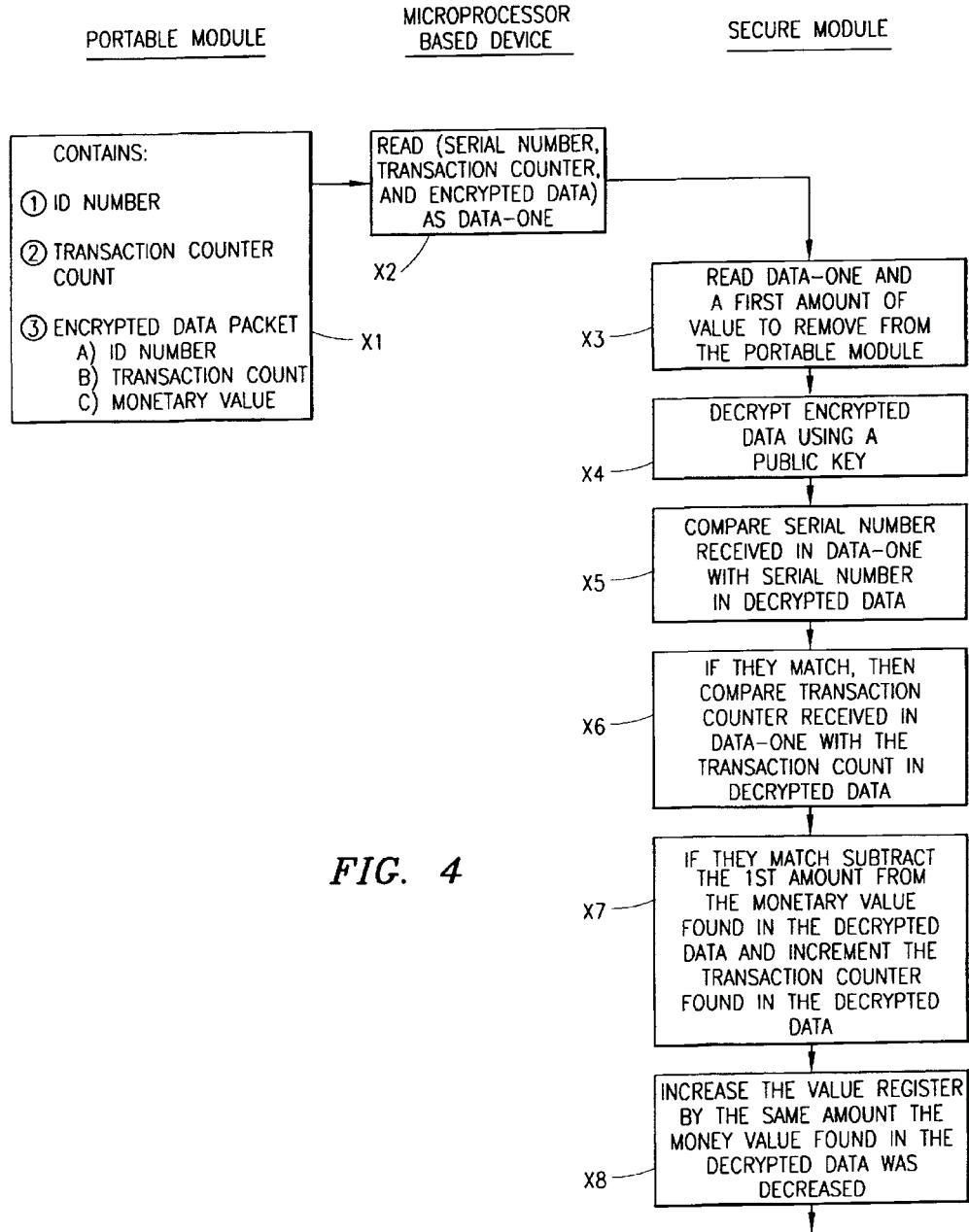
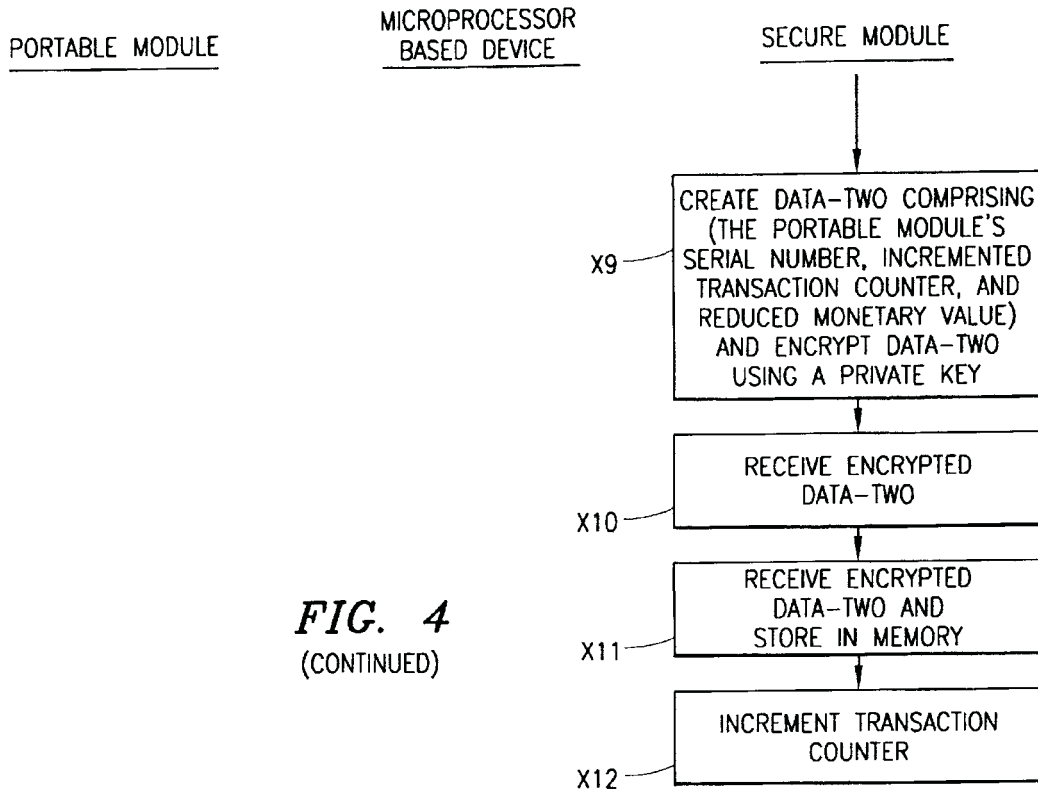


FIG. 4



**FIG. 4**  
(CONTINUED)

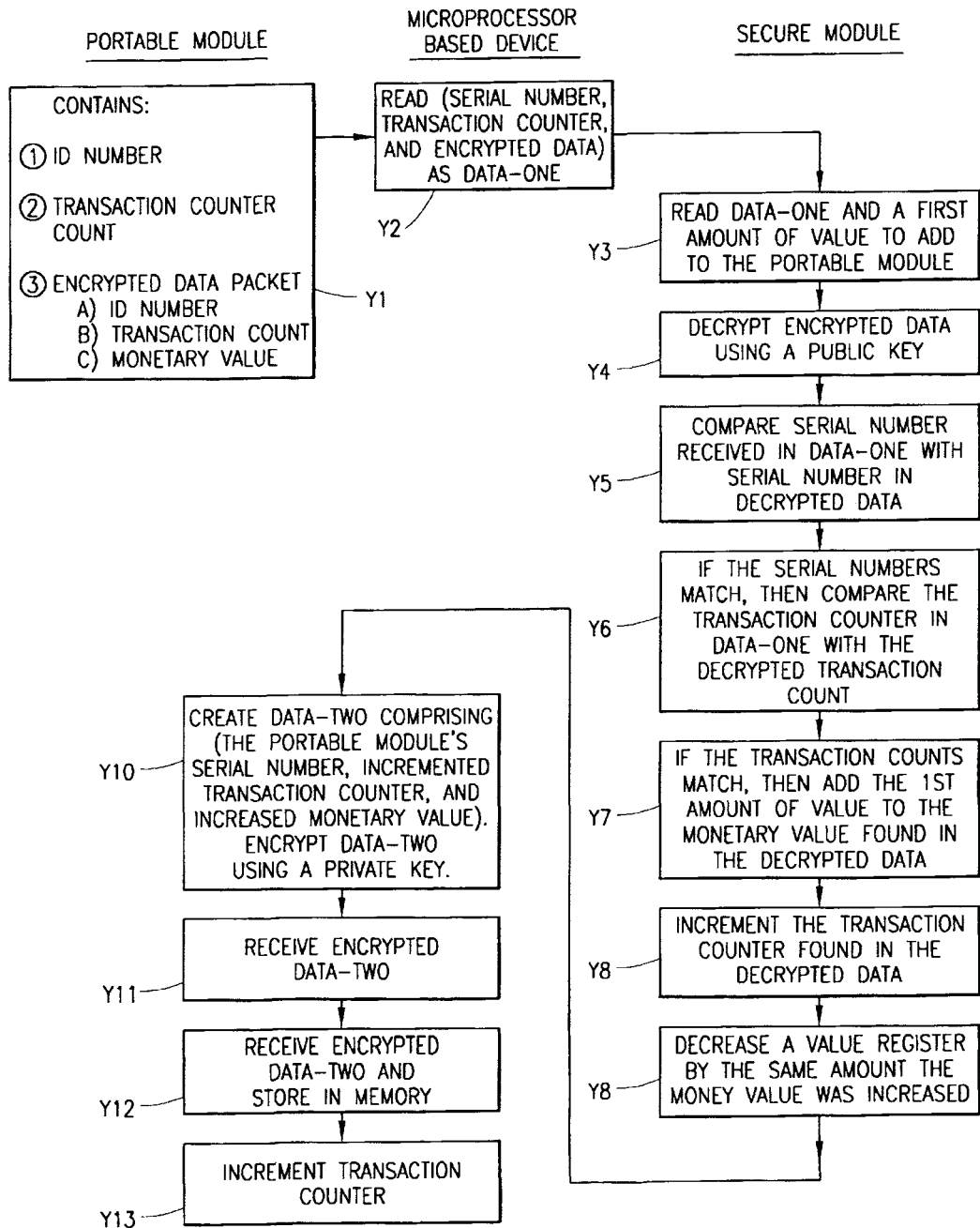


FIG. 5

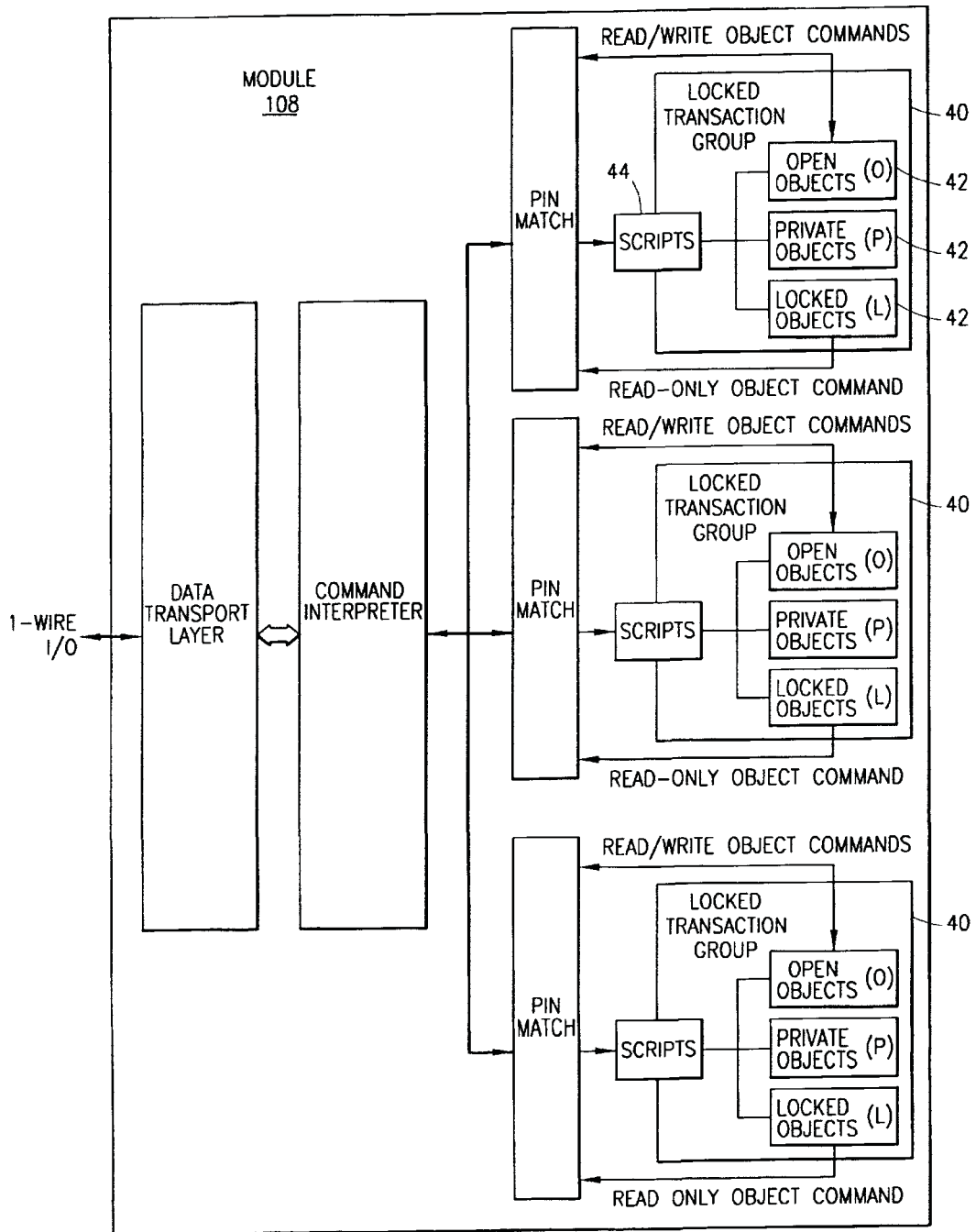


FIG. 6

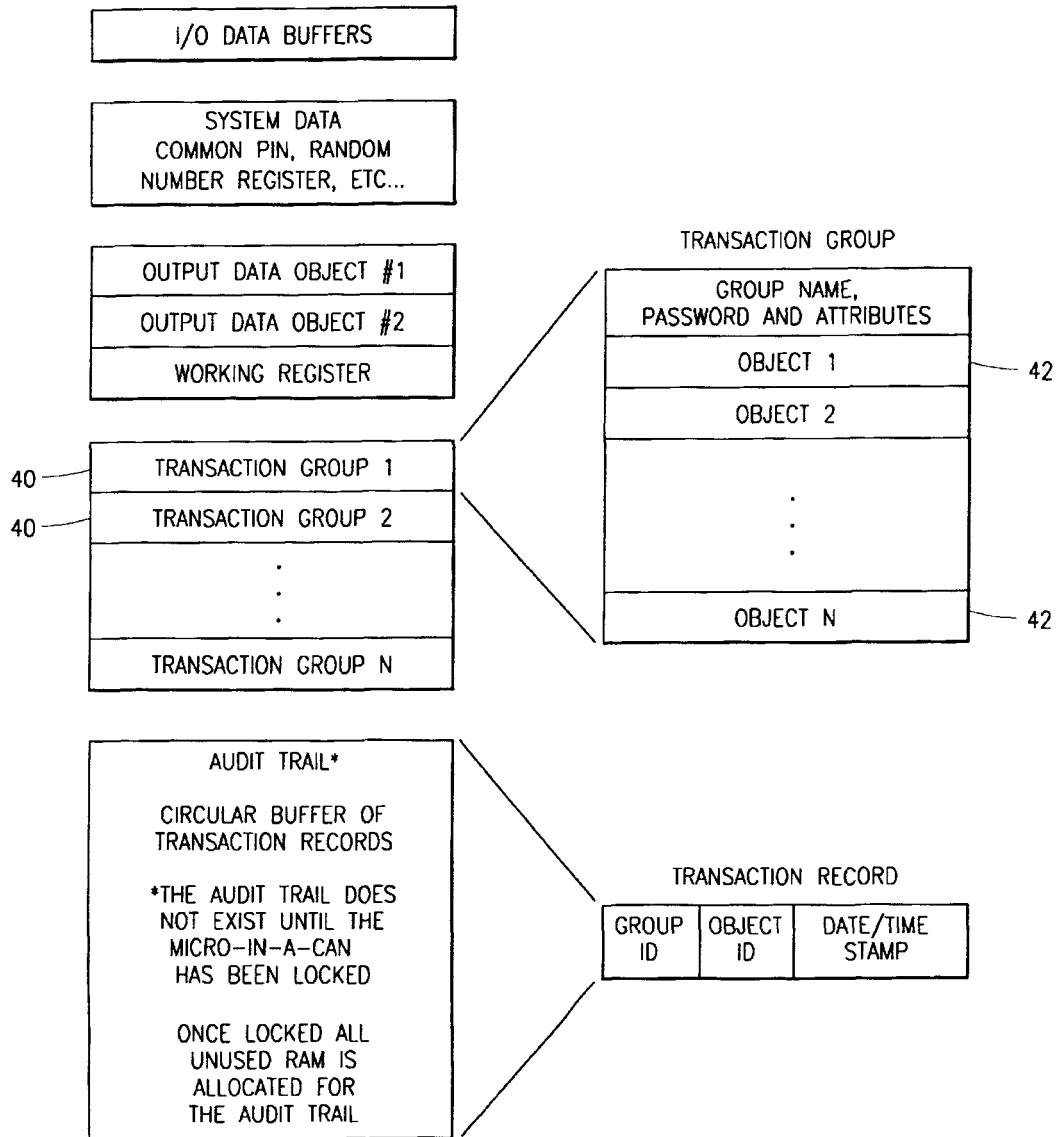


FIG. 7



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## TRANSFER OF VALUABLE INFORMATION BETWEEN A SECURE MODULE AND ANOTHER MODULE

### CROSS REFERENCE TO OTHER APPLICATIONS

The following applications of common assignee contains related subject matter and is hereby incorporated by reference:

Ser. No. 08/594,983, filed Jan. 31, 1996, entitled METHOD, APPARATUS, SYSTEM AND FIRMWARE FOR SECURE TRANSACTIONS; and

Ser. No. 08/595,014, filed Jan. 31, 1996, entitled METHOD, APPARATUS AND SYSTEM FOR TRANSFERRING UNITS OF VALUE.

### BACKGROUND OF THE INVENTION

#### 1. Technical Field of the Invention

The present invention relates to a method and system for transferring valuable information securely between a secure module and another module. More particularly, the present invention relates to transferring units of value between a microprocessor based secure module and another module used for carrying a monetary equivalent.

#### 2. Description of Related Art

In the past the preferred means for paying for an item was cash. As our society has become more advanced, credit cards have become an accepted way to pay for merchandise or services. The payment is not a payment to the merchant, but instead is a credit given by a bank to the user that the merchant accepts as payment. The merchant collects money from the bank based on the credit. As time goes on, cash is used less and less, and money transfers between parties are becoming purely electronic.

Present credit cards have magnetic strips to identify the owner of the card and the credit provider. Some credit cards have electronic circuitry installed that identifies the credit card owner and the credit or service provider (the bank).

The magnetic strips installed in present credit cards do not enable the card to be used as cash. That is the modern credit card does not allow the consumer to buy something with the credit card and the merchant to receive cash at the time of the transaction. Instead, when the consumer buys something on credit, the merchant must later request that the bank pay for the item that the consumer bought. The bank then bills the consumer for the item that was bought.

Thus, there is a need for an electronic system that allows a consumer to fill an electronic module with a cash equivalent in the same way a consumer fills his wallet with cash. When the consumer buys a product or service from a merchant, the consumer's module can be debited and the merchant's cash drawer can be credited without any further transactions with a bank or service provider.

### SUMMARY OF THE INVENTION

The present invention is an apparatus, system and method for communicating a cash equivalent electronically to and from a portable module. The portable module can be used as a cash equivalent when buying products and services in the market place.

The present invention comprises a portable module that can communicate to a secure module via a microprocessor based device. The portable module can be carried by a consumer, filled with electronic money at an add-money

2

station, and be debited by a merchant when a product or service is purchased by the consumer. As a result of a purchase, the merchant's cash drawer will indicate an increase in cash value.

### BRIEF DESCRIPTION OF THE DRAWINGS

A more complete understanding of the method and apparatus of the present invention may be had by reference to the following Detailed Description when taken in conjunction with the accompanying Drawings wherein:

FIG. 1 depicts an exemplary system for transferring valuable information between a module and a secure device;

FIG. 2 is a block diagram of an embodiment of a portable module;

FIG. 3 is a block diagram of an embodiment of a microprocessor based module;

FIG. 4 is an exemplary technique for transferring valuable data securely into a portable module;

FIG. 5 is an exemplary technique for transferring valuable data securely out of a portable module;

FIG. 6 is an exemplary organization of the software and firmware within a secure microprocessor based device; and

FIG. 7 is an exemplary configuration of software and firmware within a secure microprocessor based device.

### DETAILED DESCRIPTION OF A PRESENTLY PREFERRED EXEMPLARY EMBODIMENT

FIG. 1 depicts a block diagram of an exemplary system **100** for transferring valuable information to and from a portable module. A portable module **102**, which will be described in more detail later, communicates to a microprocessor based device **104**. The portable module **102** may contain information that represents units of exchange or a currency equivalent. The microprocessor based device **104** can be any of an unlimited number of devices. For example, the microprocessor based device **104** could be a personal computer, an add-a-fare machine at a train or bus station (similar to those in today's District of Columbia metro stations), a turn style, a toll booth, a bank's terminal, a ride at a carnival, a washing machine at a laundromat, a locking device, a mail metering device or any device that controls access, or meters a monetary equivalent, etc.

The means for communication **106** between the portable module **102** and the microprocessor based device **104** is preferably via a single wire or contact connection. The single wire connection **106** preferably incorporates a communication protocol that allows the portable module **102** and the microprocessor based device **104** to communicate in a bidirectional manner. Preferably the communication protocol is a one-wire protocol developed by Dallas Semiconductor. It is understood that the means for communicating **106** is not limited to a single wire connection. The communication means **106** could be multiple wires, a wireless communication system, infrared light, any electromagnetic means, a magnetic technique, or any other similar technique.

The microprocessor based device **104** is electrically connected to another microprocessor based device, which is preferably a secure device **108**. The term secure device means that the device is designed to contain a secret code and the secret code is extremely difficult to learn. An example of a secure device **108** is explained later in this document.

The microprocessor based device **104** can be connected to a variety of other devices. Such devices include, but are not

limited to a cash acceptor **110**, an automatic teller machine (ATM) **112**, a credit card reader **114**, and a phone line **116**.

The cash acceptor **110** is adapted to receive cash in the form of currency, such as dollar bills or coins. The cash acceptor **110**, preferably, determines the value of the accepted currency. The cash acceptor **110** communicates to the microprocessor based device **104** and informs the device **104** of how much currency has been deposited in the cash acceptor **110**.

The cash acceptor **110** can also be a device which provides currency. That is, the cash acceptor **110** in response to a communication from the microprocessor based device **104**, may provide a metered amount of currency to a person.

The credit card reader **114**, and ATM **112** can also be attached to the microprocessor based device **104**. The credit card reader **114** could be used to read a user's credit card and then, when authorized, either communicate to the microprocessor based device **104** that units of exchange need to be added to the portable module or that units of exchange need to be extracted from the portable module to pay for a good, service or credit card bill.

The ATM **112** may also be connected to the microprocessor based device. Via communications from the ATM **112**, the microprocessor based device **104** can be informed that units of exchange need to be added or subtracted from the portable module **102**.

Furthermore, it is also possible that the microprocessor based device **104** is connected to a phone line **116**. The phone line may be used for a variety of things. Most importantly, the phone line may be used to allow the microprocessor based device **104** to communicate with a network of devices. Such telephonic communication may be for validating transactions or for aiding the accounting of transactions that are performed via the microprocessor based device's **104** aid. It is further understood that the phone line may be any of a vast variety of communication lines including wireless lines. Video, analog, or digital information may be communicated over the phone line **116**.

FIG. 2 depicts a preferred exemplary portable module **102**. The portable module **102** is preferably a rugged read/write data carrier that can act as a localized data base and be easily accessed with minimal hardware. The module can be incorporated in a vast variety of portable items which includes, but is not limited to a durable micro-can package that is highly resistant to environmental hazards such as dirt, moisture, and shock. The module can be incorporated into any object that can be articulated by a human or thing, such as a ring, bracelet, wallet, name tag, necklace, baggage, machine, robotic device, etc. Furthermore, the module **102** could be attached to a stationary item and the microprocessor based device **104** may be articulated to the portable module **102**. For example, the module **102** may be attached to a piece of cargo and a module reader may be touched to or brought near the module **102**. The module reader may be part of the microprocessor based device **104**.

The portable module **102** comprises a memory **202** that is preferably, at least in part, nonvolatile memory for storing and retrieving vital information pertaining to the system to which the module **102** may become attached to. The memory **202** may contain a scratchpad memory which may act as a buffer when writing into memory. Data is first written to the scratchpad where it can be read back. After data has been verified, the data is transferred into the memory.

The module **102** also comprises a counter **206** for keeping track of the number of transactions the module has per-

formed (the number of times certain data in the memory of the module has been changed). A timer **102** may be provided in the module to provide the ability to time stamp transactions performed by the module. A memory controller **204** controls the reading and writing of data into and out of the memory **202**.

The module also may comprise an identification number **210**. The identification number preferably uniquely identifies the portable module from any other portable module.

An input/output control circuit **212** controls the data flow into and out of the portable module **102**. The input/output control ("I/O") **212** preferably has an input buffer and an output buffer and interface circuitry **214**. As stated above, the interface circuitry **214** is preferably a one-wire interface. Again, it is understood that a variety of technologies can be used to interface the portable module **102** to another electronic device. A single wire or single connection is preferred because the mechanics of making a complete connection is simplified. It is envisioned that a proximity/wireless communication technique is also a technique for communicating between the module **102** and another device. Thus, the interface circuit **214** can be a single wire, multiple wire, wireless, electromagnetic, magnetic, light, or proximity, interface circuit.

FIG. 3 depicts a block diagram of an exemplary secure microprocessor based device ("secure device") **108**. The secure device circuitry can be a single integrated circuit. It is understood that the secure device **108** could also be a monolithic or multiple circuits combined together. The secure device **108** preferably comprises a microprocessor **12**, a real time clock **14**, control circuitry **16**, a math coprocessor **18**, memory circuitry **20**, input/output circuitry **26**, and an energy circuit **34**.

The secure device **108** could be made small enough to be incorporated into a variety of objects including, but not limited to a token, a card, a ring, a computer, a wallet, a key fob, a badge, jewelry, a stamp, or practically any object that can be grasped and/or articulated by a user of the object. In the present system **100**, the secure device **108** is preferably adapted to be a trusted certifying authority. That is the secure device **108** is a trusted computer. The secure device **108** comprises a numeric coprocessor **18** optimized for math intensive encryption. The BIOS is immune to alteration and is specifically designed for secure transactions. This secure device **108** is preferably encased in a durable, dirt, moisture and shock resistant stainless steel enclosure, but could be encased in wide variety of structures so long as specific contents of the secure device **108** are extremely difficult to decipher. The secure device **108**. The secure device **108** may have the ability to store or create a private/public key set, whereby the private key never leaves the secure device **108** and is not revealed under almost any circumstance. Furthermore, the secure module **108** is preferably designed to prevent discovery of the private key by an active self-destruction of the key upon wrongful entry.

The microprocessor **12** is preferably an 8-bit microprocessor, but could be 16, 32, 64 or any operable number of bits. The clock **14** provides timing for the module circuitry. There can also be separate clock circuitry **14** that provides a continuously running real time clock.

The math coprocessor circuitry **18** is designed and used to handle very large numbers. In particular, the coprocessor will handle the complex mathematics of RSA encryption and decryption or other types of math intensive encryption or decryption techniques.

The memory circuitry **20** may contain both read-only-memory and non-volatile random-access-memory.

Furthermore, one of ordinary skill in the art would understand that volatile memory, EPROM, SRAM and a variety of other types of memory circuitry might be used to create an equivalent device.

Control circuitry 16 provides timing, latching and various necessary control functions for the entire circuit.

An input/output circuit 26 enables bidirectional communication with the secure module 108. The input/output circuitry 26 preferably comprises at least an output buffer and an input buffer. For communication via a one-wire bus, one-wire interface circuitry can be included with the input/output circuitry 26. It is understood that the input/output circuitry 26 of the secure device 108 can be designed to operate on a single wire, a plurality of wires or any means for communicating information between the secure module 108 and the microprocessor based device 104.

An energy circuit 34 may be necessary to maintain stored information in the memory circuitry 20 and/or aid in powering the other circuitry in the module 108. The energy circuit 34 could consist of a battery, capacitor, R/C circuit, photo-voltaic cell, or any other equivalent energy producing circuit or means.

The firmware architecture of the secure module 108 and how it operates within the exemplary system for transferring valuable information, such as units of exchange or currency, between the secure module 108 and a portable module 102 will now be discussed. The secure module 108 provides encryption and decryption services for confidential data transfer through the microprocessor based device 104. The following examples are intended to illustrate a preferred feature set of the secure module 108 and to explain the services that the exemplary system 100 can offer. These applications and examples by no means limit the capabilities of the invention, but instead bring to light a sampling of its capabilities.

#### I. OVERVIEW OF THE PREFERRED SECURE MODULE 108 AND ITS FIRMWARE DESIGN

Referring to FIG. 3 again, the secure module 108 preferably contains a general-purpose, 8051-compatible micro controller 12 or a reasonably similar product, a continuously running real-time clock 14, a high-speed modular exponentiation accelerator for large integers (math coprocessor) 18, input and output buffers 28, 30 with a one-wire interface 32 for sending and receiving data, 32 Kbytes of ROM memory 22 with preprogrammed firmware, 8 Kbytes of NVRAM (non-volatile RAM) 24 for storage of critical data, and control circuitry 16 that enables the micro controller 12 to be powered up to interpret and act on the data placed in an input data object. The module 108 draws its operating power from a single wire, one-wire communication line. The micro controller 12, clock 14, memory 20, buffers 28, 30, one-wire front-end 32, modular exponentiation accelerator 18, and control circuitry 16 are preferably integrated on a single silicon chip and packaged in a stainless steel micro can using packaging techniques which make it virtually impossible to probe the data in the NVRAM 24 without destroying the data. Initially, most of the NVRAM 24 is available for use to support applications such as those described below. One of ordinary skill will understand that there are many comparable variations of the module design. For example, volatile memory might be used, or an interface other than a one-wire interface could be used.

The secure module 108 is preferably intended to be used first by a Service Provider who loads the secure module 108 with data to enable it to perform useful functions, and second by an End User who issues commands to the secure module 108 to perform operations on behalf of the Service

Provider for the benefit of the End User. For this reason, the secure module 108 offers functions to support the Service Provider in setting up the module for an intended application. It also offers functions to allow the End User to invoke the services offered by the Service Provider.

Each Service Provider can reserve a block of NVRAM memory to support its services by creating a transaction group 40 (refer to FIGS. 6 and 7). A transaction group 40 is simply a set of software objects 42 that are defined by the Service Provider. These objects 42 include both data objects (encryption keys, transaction counts, money amounts, date/time stamps, etc.) and transaction scripts 44 which specify how to combine the data objects in useful ways. Each Service Provider creates his own transaction group 40, which is independent of every other transaction group 40. Hence, multiple Service Providers can offer different services in the same module 108. The number of independent Service Providers that can be supported depends on the number and complexity of the objects 42 defined in each transaction group 40. Examples of some of the objects 42 that can be defined within a transaction group 40 are the following:

RSA Modulus	Clock Offset
RSA Exponent	Random SALT
Transaction Script	Configuration Data
Transaction Counter	Input Data
Money Register	Output Data
Destructor	

Within each transaction group 40 the secure module 108 will initially accept certain commands which have an irreversible effect. Once any of these irreversible commands are executed in a transaction group 40, they remain in effect until the end of the module's useful life or until the transaction group 40, to which it applies, is deleted from the secure module 108. In addition, there are certain commands which have an irreversible effect until the end of the module's life or until a master erase command is issued to erase the entire contents of the secure module 108. These commands will be discussed further below. These commands are essential to give the Service Provider the necessary control over the operations that can be performed by the End User. Examples of some of the irreversible commands are:

Privatize Object	Lock Object
Lock Transaction Group	Lock Micro-In-A-Can™

Since much of the module's utility centers on its ability to keep a secret, the Privatize command is a very important irreversible command.

Once the secure module 108, as a whole, is locked, the remaining NVRAM memory 24 is allocated for a circular buffer for holding an audit trail of previous transactions. Each of the transactions are identified by the number of the transaction group, the number of objects 42 within the specified group, and the date/time stamp.

The fundamental concept implemented by the firmware is that the Service Provider can store transaction scripts 44 in a transaction group 40 to perform only those operations among objects that he wishes the End User to be able to perform. The Service Provider can also store and privatize RSA key or keys (encryption keys) that allow the secure module 108 to "sign" transactions on behalf of the Service Provider, thereby guaranteeing their authenticity. By privatizing and/or locking one or more objects 42 in the trans-

action group 40, the Service Provider maintains control over what the secure module 108 is allowed to do on his behalf. The End User cannot add new transaction scripts 44 and is therefore limited to the operations on objects 42 that can be performed with the transaction scripts 44 programmed by the Service Provider.

## II. USAGE MODELS OF THE SECURE MODULE 108 AND PORTABLE MODULE 102

This section presents practical applications of the system 100. Each of these applications is described in enough detail to make it clear why the secure module 108 and portable module 102 are important to the system application.

### A. TRANSFERRING UNITS OF EXCHANGE OUT OF A PORTABLE MODULE 102

This section describes an example of how a portable module 102 and a secure module 108 operate in conjunction with the microprocessor based device 104 so that units of exchange can be securely transferred out of the portable module 102 and deposited into the secure module 108 and/or potentially communicated to at least one of the cash acceptor 110, ATM 112, credit card reader 114, or the phone line 116.

Referring to FIG. 4, initially the portable module 102 contains its ID number, a count within its transaction counter and an encrypted data packet stored in memory. Encrypted within the data packet is the portable modules ID number, the portable modules transaction count number, and the amount of value (the monetary value) of the portable module at the present time X1.

The user of the portable module touches, or somehow puts the portable module 102 into communication with the microprocessor based device 104. For explanation purposes, suppose the portable module 102 is being used as a token used to pay for a train fare. Thus, the microprocessor based device 104 could be, in this case, a turn style that allows the user to enter a train platform. The cost of entering the train platform is known by the microprocessor based device 104.

The microprocessor based device 104 reads the portable module's serial number, transaction count, and the encrypted data packet X2. This data could be referred to as a first data.

The microprocessor device 104 then provides the first data along with a first value, being the amount of value to be debited from the portable token (the train fare), to the secure module 108 X3. The secure module 108 decrypts the encrypted data found in the first data using a public key X4.

Next, the secure module 108 makes a few comparisons to make sure that the data received is good data and not counterfeit. The secure module 108 compares the serial number received in the first data with the decrypted serial number X5. If the two serial numbers match then the secure module 108 compares the transaction count received in the first data with the decrypted transaction count X6. If the two transaction counts match then the secure module is comfortable that the data received is not counterfeit data. It is understood that the comparisons can be done in any order.

Furthermore, there may have been a time stamp sent from the portable module 102. The time stamp may indicate a variety of things. One thing could be an indication of whether the portable module is still valid or the time stamp may further enable the secure module to decide if the data is or is not counterfeit.

Assuming all the data passed to the secure module 108 is determined to be valid data, the secure module 108 subtracts the first value, the train fare, from the monetary value of the portable module 102 X7. The decrypted transaction count is then incremented.

A register within the secure module 108 is increased by the amount of the first value, the train fare, so that the secure

module can keep an accounting of the amount of "money" it has collected X8. The secure module 108 creates a data packet, a second data, which comprises at least the portable module's serial number, the incremented transaction count, and the reduced monetary value of the portable module 102. The second data packet is then encrypted by the secure module 108 using a private key X9.

The microprocessor based device 104 receives the encrypted second data packet, passes the encrypted second data packet to the portable module 102 X10, and opens the turn style to let the module's user onto the train platform. The portable module 102 receives the encrypted second data packet and stores it in memory X11. The portable module also increments its transaction count indicating that another transaction has occurred X12.

Thus, the above description indicates how valuable information can be transferred between a portable insecure module 102 and a secure module 108 wherein there is a conservation of value. That is, no value is gained or lost. Value that was in the portable module 102 was decreased by the same amount value was added to the secure module 108. In the example provided, the decrease and increase in value was equal to a train fare. Such an increment or decrement can also be equal to an amount provided by an ATM, credit card transaction, cash acceptor, etc.

It is also understood that the insecure portable module 102 could be another secure module similar to the secure module in the system, but programed to act like a portable module 102.

### B. TRANSFERRING UNITS OF EXCHANGE INTO THE PORTABLE MODULE 102

In this example, for simplicity, suppose the portable module does not have any monetary value and the user of the portable module wishes to "fill it up" with value. Suppose the user wishes to take cash out of an ATM machine and instead of pocketing the cash, the user wishes to put the cash value into the portable module 102.

Referring to FIG. 5, the portable module 102 contains its ID number, a transaction count and an encrypted data packet containing the portable module's ID number, transaction count and the monetary value of the portable module 102 Y1. The microprocessor based device 104, which in this example could be part of the ATM machine 112, receives the information contained in the portable module 102 when a communication is initiated between the portable module 102 and the microprocessor based device 104 Y2.

The microprocessor based device 104 passes the module's serial number, transaction count, and encrypted data packet as a first data packet to the secure module 108. The microprocessor based device also passes the amount of amount of monetary value to add to the portable module 102, as indicated by the ATM 112, to the secure module 108 Y3.

The secure module 108 decrypts the encrypted data passed to it using a public key Y4. The secure module 108 then makes a few comparisons to make sure that the data it has just received is valid and not counterfeit. The secure module 108 compares the serial number (ID number) received in the first data packet with the serial number (ID number) found in the decrypted data Y5. The secure module 108 also compares the transaction count passed the first data packet with the transaction count found in the decrypted data Y6. If the serial numbers and transaction counters match, then the secure module decides that the data received is valid and the secure module adds the monetary value, indicated by the ATM to the monetary value of the decrypted data Y7. The decrypted transaction count is incremented Y8. A register within the secure module may be decremented by the

same amount that the monetary value of the decrypted data was increased Y8.

The secure module 108 creates a second data packet, that contains the portable module's ID number, the incremented transaction counter and the increased monetary value. The second data packet is then encrypted using a private key Y10.

The microprocessor based device 104 reads the encrypted second data packet and sends it to the portable module 102 Y11. The portable module receives the encrypted second data packet and stores it in memory Y12. The portable module also advances its transaction counter Y13. The result being that the portable module now has the value of the cash withdrawn from the ATM 112. Furthermore, a record of the transaction may have been recorded and kept in the secure module, as well as by the bank that operates the ATM 112.

Exemplary Firmware Definitions for Use With the Secure Module		
Object	The most primitive data structure accepted by and operated on by the secure modules firmware. A list of valid objects and their definitions is provided in the next section.	20
Group	A self-contained collection of objects. An object's scope is restricted to the group of which it is a member.	25
Group ID	A number preferably between 0 and 255 representing a specific group.	
Object ID	A number preferably between 0 and 255 representing a specific object within a specific group.	30
Object Type	Preferably a 1-byte type specifier that describes a specific object.	
PIN	An alphanumeric Personal Identification number that is preferably eight bytes in length.	
Common PIN	The PIN that controls access to shared resources such as the audit trail. It is also used to control the host's ability to create and delete groups.	35
Group PIN	The PIN that controls access to operations specific to objects within a group.	40
Audit Trail	A record of transactions occurring after the secure module has been locked.	
Locked Object	An object which has been locked by executing the lock object command. Once an object is locked it is not directly readable.	45
Private Object	An object which has been privatized by executing the privatize object command. Once an object is private, it is not directly readable or writable.	50
Locked Group	A group which has been locked using the locked group command. After a group has been locked it will not allow object creation.	
Composite Object	A combination of several objects. The individual objects inherit the attributes of the composite object.	55

Exemplary Object Definitions		
RSA Modulus	A large integer preferably of at most 1024 bits in length. It is the product of 2 large prime numbers that are each about half the number of bits in length of the desired modulus size. The RSA modulus is	65

Exemplary Object Definitions		
	used in the following equations for encrypting and decrypting a message M: Encryption: $C = M^e \pmod{N}$ (1) Decryption: $M = C^d \pmod{N}$ (2) where C is the cyphertext, d and e are the RSA exponents (see below) and N is the RSA modulus. Both e and d (shown in equations 1 and 2 above) are RSA exponents. They are typically large numbers but are smaller than the modulus (N). RSA exponents can be either private or public. When RSA exponents are created in the secure module, they may be declared as either. Once created an exponent may be changed from a public exponent to a private exponent. After an exponent has been made private, however, it will remain private until the transaction group 40 to which it belongs is destroyed.	
RSA Exponent		
Transaction Script	A transaction script is a series of instructions to be carried out by the secure module. When invoked the secure module firmware interprets the instructions in the script and places the results in the output data object (see below). The actual script is simply a list of objects. The order in which the objects are listed specifies the operations to be performed on the objects. transaction scripts 44 preferably may be as long as 128 bytes.	
Transaction Counter	The transaction counter object is preferably 4 bytes in length and is usually initialized to zero when it is created. Every time a transaction script, which references this object, is invoked, the transaction counter increments by 1. Once a transaction counter has been locked it is read only and provides an irreversible counter.	
Money Register	The money register object is preferably 4 bytes in length and may be used to represent money or some other form of credit. Once this object has been created, it must be locked to prevent a user from tampering with its value. Once locked the value of this object can be altered only by invoking a transaction script. A typical transaction group 40 which performs monetary transactions might have one script for withdrawals from the money register and one for deposits to the money register.	
Clock Offset	This object is preferably a 4 byte number which contains the difference between the reading of the secure module's real-time clock and some convenient time (e.g., 12:00 a.m., January 1, 1970). The true time can then be obtained from the secure module by adding the value of the clock offset to the real-time clock.	
SALT	A SALT object is preferably 20 bytes in length and should be initialized with random data when it is created. When a host transmits a generate random SALT command, the secure module combines the previous SALT with the secure module's random number (produced preferably by	

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Exemplary Object Definitions	
Configuration Data	randomly occurring power-ups) to generate a new random SALT. If the SALT object has not been privatized it may subsequently be read by issuing a read object command. This is a user defined structure with preferably a maximum length of 128 bytes. This object is typically used to store configuration information specific to its transaction group 40. For example, the configuration data object may be used to specify the format of the money register object (i.e., the type of currency it represents) Since this object has no pre-defined structure, it may never be used by a transaction object.
Input Data	An input data object is simply an input buffer with preferably a maximum length of 128 bytes. A transaction group may have multiple input objects. The host uses input data objects to store data to be processed by transaction scripts 44.
Output Data	The output data object is used by transaction scripts as an output buffer. This object is automatically created when the transaction group is created. It is preferably 512 bytes in length and inherits password protection from its group.
Random Fill	When the script interpreter encounters this type of object it automatically pads the current message so that its length is 1 bit smaller than the length of the preceding modulus. A handle to this object is automatically created when the transaction group is created. It is a private object and may not be read using the read object command.
Working Register	This object is used by the script interpreter as working space and may be used in a transaction script. A handle to this object is automatically created when the transaction group is created. It is a private object and may not be read using the read object command.
ROM Data	This object is automatically created when the transaction group is created. It is a locked object and may not be altered using the write object command. This object is 8 bytes and length and its contents are identical to the 8 by ROM data of the Micro-In-A-Can™.
PIN_TO_ERASE	00000001b (require PIN for Master Erase)
PIN_TO_CREATE	00000010b (require PIN for group creation)

Initially the secure module has a PIN (Personal Identification Number) of 0 (Null) and an option byte of 0. Once a PIN has been established it can only be changed by providing the old PIN or by a Master Erase. However, if the

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PIN\_TO\_ERASE bit is set in the option byte, the PIN can only be changed through the set common PIN command. Possible error codes for the set common PIN command:

5	ERR_BAD_COMMON_PIN	(Common PIN match failed)
	ERR_BAD_PIN_LENGTH	(New PIN length > 8 bytes)
10	ERR_BAD_OPTION_BYTE	(Unrecognizable option byte)

For all commands described in this section, data received by the host will be in the form of a return packet. A return packet has the following structure:

15	Command status byte (0 if command successful, error code otherwise, 1 byte)	
	Output data length	(Command output length, 2 bytes)
20	Output data	(Command output, length specified above).

Master Erase (02H)

25	Transmit data	02H, Common PIN
	Receive data	CSB=0 if command was successful, ERR_BAD_COMMON_PIN otherwise
30	Output length=0	Output data=0

Notes:

If the LSB (least significant bit) of the PIN option is clear (i.e. PIN not required for Master Erase) then a 0 is transmitted for the Common PIN value. In general this text will always assume a PIN is required. If no PIN has been established a 0 should be transmitted as the PIN. This is true of the common PIN and group PINS (see below). If the PIN was correct the firmware deletes all groups (see below) and all objects within the groups. The common PIN and common PIN option byte are both reset to zero.

After everything has been erased the secure module transmits the return packet. The CSB is as described above. The output data length and output data fields are both set to 0.

Create Group (03H)

40	Transmit data	03H, Common PIN, Group name, Group PIN
50	Receive data	CSB=0 if command successful, appropriate error code otherwise
	Output length=1 if successful, 0 otherwise	Output data=Group ID if successful, 0 otherwise

Notes:

The maximum group name length is 16 bytes and the maximum PIN length is eight bytes. If the PIN\_TO\_CREATE bit is set in the common PIN option byte and the PIN transmitted does not match the common PIN the secure module will set the OSC to ERR\_BAD\_COMMON\_PIN.

Possible error return codes for the create group command:

65	ERR_BAD_COMMON_PIN	(Incorrect common PIN)
	ERR_BAD_NAME_LENGTH	(If group name length > 16 bytes)

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ERR_BAD_PIN_LENGTH	(If group PIN length > 8 bytes)	
ERR_MIAC_LOCKED	(The secure module has been locked)	5
ERR_INSUFFICIENT_RAM	(Not enough memory for new group)	

Set Group PIN (04H)

Transmit data 10  
 04H, Group ID, old GPIN, new GPIN  
 Receive data  
 CSB=0 if command successful, appropriate error code otherwise  
 Output length=0  
 Output data=0  
 Notes:  
 The Group PIN only restricts access to objects within the group specified by the group ID transmitted in the command packet.  
 Possible error codes for the set group PIN command:

ERR_BAD_GROUP_PIN	(Group PIN match failed)	
ERR_BAD_PIN_LENGTH	(New group PIN length > 8 bytes)	

Create Object (05H)

Transmit data 30  
 05H, Group ID, Group PIN, Object type, Object attributes, Object data  
 Receive data  
 CSB=0 if command successful, appropriate error code otherwise 35  
 Output length=1 if successful, 0 otherwise  
 Output data=object ID if successful, 0 otherwise  
 Notes:  
 If the Create Object command is successful the secure module firmware returns the object's ID within the group specified by the Group ID. If the PIN supplied by the host was incorrect or the group has been locked by the Lock Group command (described below) the secure module returns an error code in the CSB. An object creation will also fail if the object is invalid for any reason. For example, if the object being created is an RSA modulus (type 0) and it is greater than 1024 bits in length. transaction script creation will succeed if it obeys all transaction scripts rules.  
 Possible error return codes for the create object command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)	
ERR_GROUP_LOCKED	(The group has been locked)	
ERR_MIAC_LOCKED	(The secure module has been locked)	55
ERR_INVALID_TYPE	(The object type specified is invalid)	
ERR_BAD_SIZE	(The objects length was invalid)	
ERR_INSUFFICIENT_RAM	(Not enough memory for new object)	
<u>Object types:</u>		
RSA modulus	0	
RSA exponent	1	
Money register	2	
Transaction counter	3	

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Transaction script	4
Clock offset	5
Random SALT	6
Configuration object	7
Input data object	8
Output data object	9
<u>Object Attributes:</u>	
Locked	00000001b
Privatized	00000010b

Objects may also be locked and privatized after creation by using the Lock Object and Privatize Object commands described below.

Lock Object (06H)

Transmit data  
 06H, Group ID, Group PIN, Object ID  
 Receive data  
 CSB=0 if command successful, appropriate error code otherwise  
 Output length=0  
 Output data=0  
 Notes:  
 If the Group ID, Group PIN and Object ID are all correct, the secure module will lock the specified object. Locking an object is an irreversible operation.  
 Possible error return codes for the lock object command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)	
ERR_GROUP_LOCKED	(The group has already been locked)	
ERR_MIAC_LOCKED	(The secure module has been locked)	
ERR_BAD_GROUP_ID	(Specified group does not exist)	35
ERR_BAD_OBJECT_ID	(Specified object does not exist)	

Privatize Object (07H)

Transmit data 40  
 07H, Group ID, Group PIN, Object ID  
 Receive data  
 CSB=0 if successful, appropriate error code otherwise  
 Notes:  
 If the Group ID, Group PIN and Object ID were valid the object will be privatized. Privatized objects share all the properties of locked objects but are not readable. Privatized objects are only modifiable through transaction scripts. Note that locking a privatized object is legal, but has no meaning since object privatization is a stronger operation than object locking. Privatizing an object is an irreversible operation.  
 Possible error return codes for the privatize object command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)	
ERR_GROUP_LOCKED	(The group has already been locked)	
ERR_MIAC_LOCKED	(The secure module has been locked)	
ERR_BAD_GROUP_ID	(Specified group does not exist)	60
ERR_BAD_OBJECT_ID	(Specified object does not exist)	

Make Object Destructable (08H)

Transmit data  
 08H, Group ID, Group PIN, Object ID

15

Receive data

CSB=0 if successful, appropriate error code otherwise

Notes:

If the Group ID, Group PIN and Object ID were valid the object will be made destructable. If an object is destructable it becomes unusable by a transaction script after the groups destructor becomes active. If no destructor object exists within the transaction group the destructible object attribute bit has no affect. Making an object destructable is an irreversible operation.

Possible error return codes for the make object destructable command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_GROUP_LOCKED	(The group has already been locked)
ERR_MIAC_LOCKED	(The secure module has been locked)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_OBJECT_ID	(Specified object does not exist)

Lock Secure module (09H)

Transmit data

09H, Common PIN

Receive data

CSB=0 if successful, appropriate error code otherwise

Output length=2 if successful, 0 otherwise

Output data=audit trail size if successful, 0 otherwise

Notes:

If the host supplied Common PIN is correct and the secure module has not previously been locked, the command will succeed. When the secure module is locked it will not accept any new groups or objects. This implies that all groups are automatically locked. The RAM not used by the system or by groups will be used for an audit trail. There is no audit trail until the secure module has successfully been locked!

An audit trail record is six bytes long and has the following structure:

Group ID|Object ID|Date/Time stamp.

Once an audit trail has been established, a record of the form shown above will be stored in the first available size byte location every time a transaction script is executed. Note that since the secure module must be locked before the audit trail begins, neither the group ID nor any object ID is subject to change. This will always allow an application processing the audit trail to uniquely identify the transaction script that was executed. Once the audit trail has consumed all of its available memory, it will store new transaction records over the oldest transaction records.

Possible error codes for the lock secure module command:

ERR_BAD_COMMON_PIN	(Supplied common PIN was incorrect)
ERR_MIAC_LOCKED	(Secure module was already locked)

Lock Group (0AH)

Transmit data

0AH, Group ID, Group PIN

16

Receive data

CSB=0 if command successful, appropriate error code otherwise

Output length=0

Output data=0

Notes:

If the group PIN provided is correct the secure module BIOS will not allow further object creation within the specified group. Since groups are completely self-contained entities they may be deleted by executing the Delete Group command (described below).

Possible error return codes for the lock group command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_GROUP_LOCKED	(The group has already been locked)
ERR_MIAC_LOCKED	(The secure module has been locked)
ERR_BAD_GROUP_ID	(Specified group does not exist)

Invoke Transaction Script (0BII)

Transmit data

0BH, Group ID, Group PIN, Object ID

Receive data

CSB=0 if command successful, appropriate error code otherwise

Output length=1 if successful, 0 otherwise

Output data=estimated completion time

Notes:

The time estimate returned by the secure module is in sixteenths of a second. If an error code was returned in the CSB, the time estimate will be 0.

Possible error return codes for the execution transaction script command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_OBJECT_ID	(Script object did not exist in group)

Read Object (0CH)

Transmit data

0CH, Group ID, Group PIN, Object ID

Receive data

CSB=0 if command successful, appropriate error code otherwise

Output length=object length if successful, 0 otherwise

Output data=object data if successful, 0 otherwise

Notes:

If the Group ID, Group PIN and Object ID were correct, the secure module checks the attribute byte of the specified object. If the object has not been privatized the secure module will transmit the object data to the host. If the Group PIN was invalid or the object has been privatized the secure module will return a 0 in the output length, and data fields of the return packet.

Possible error codes for the read object command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_BAD_GROUP_ID	(Specified group does not exist)



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-continued

ERR_BAD_OBJECT_ID	(Object did not exist in group)
ERR_OBJECT_PRIVATIZED	(Object has been privatized)

Write Object (0DH)  
 Transmit data  
 0DH, Group ID, Group PIN, Object ID, Object size, Object Data  
 Receive data  
 CSB=0 if successful, appropriate error code otherwise  
 Output length=0  
 Output data=0

Notes:  
 If the Group ID, Group PIN and Object ID were correct, the secure module checks the attribute byte of the specified object. If the object has not been locked or privatized the secure module will clear the objects previous size and data and replace it with the new object data. Note that the object type and attribute byte are not affected.

Possible error codes for the write object command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_OBJECT_ID	(Object did not exist in group)
ERR_BAD_OBJECT_SIZE	(Illegal object size specified)
ERR_OBJECT_LOCKED	(Object has been locked)
ERR_OBJECT_PRIVATIZED	(Object has been privatized)

Read Group Name (0EH)  
 Transmit data  
 0EH, Group ID  
 Receive data  
 CSB=0  
 Output Length=length of group name  
 Output data=group name

Notes:  
 The group name length is a maximum of 16 bytes. All byte values are legal in a group name.

Delete Group (0FH)  
 Transmit data  
 0FH, Group ID, Group PIN  
 Receive data  
 CSB=0 if successful, appropriate error code otherwise  
 Output length=0  
 Output data=0

Notes:  
 If the group PIN and group ID are correct the secure module will delete the specified group. Deleting a group causes the automatic destruction of all objects within the group. If the secure module has been locked the Delete Group command will fail.

Possible error codes for the delete group command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_BAD_GROUP_ID	(Specified group does not exist)

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-continued

ERR_MIAC_LOCKED	(Secure module has been locked)
-----------------	---------------------------------

5 Get Command Status Info (10H)  
 Transmit data  
 10H  
 Receive data  
 CSB=0  
 Output length=6  
 Output data=secure module status structure (see below)

Notes:  
 15 This operation requires no PIN and never fails. The status structure is defined as follows:

Last command executed	(1 byte)
Last command status	(1 byte)
Time command received	(4 bytes)

20 Get Secure module Configuration Info (11H)  
 Transmit data  
 11H  
 Receive data  
 CSB=0  
 Output length=4  
 Output data=secure module configuration structure

Notes:  
 30 This operation requires no PIN and never fails. The configuration structure is defined as follows:

Number of groups	(1 byte)
Flag byte (see below)	(1 byte)
Audit trail size/Free RAM	(2 bytes)

The flag byte is the bitwise-or of any of the following values:

- 40 0000001b (Secure module is locked)
- 0000010b (Common PIN required for access)

Read Audit Trail Info (12H)  
 Transmit data  
 12H, Common PIN  
 Receive data  
 CSB=0 if command successful, appropriate error code otherwise  
 Output length=audit trail structure size (5) if successful, 0 otherwise  
 Output data=audit trail info structure if successful, 0 otherwise

Notes:  
 55 If the transmitted Common PIN is valid and the secure module has been locked, it returns audit trail configuration information as follows:

- Number of used transaction records (2 bytes)
  - Number of free transaction records (2 bytes)
  - 60 A boolean specifying whether or (1 byte) not the audit trail rolled since previous read command
- Possible error codes for the read audit trail info command:  
 ERR\_BAD\_COMMON\_PIN (Common PIN was incorrect)  
 ERR\_MIAC\_NOT\_LOCKED (Secure module is not locked)

Read Audit Trail (13H)

Transmit data  
 13H, Common PIN  
 Receive data  
 CSB=0 if command successful, appropriate error code otherwise  
 Output length=# of new records \* 6 if successful, 0 otherwise  
 Output data=new audit trail records

Notes:

If the transmitted common PIN is valid and the secure module has been locked, it will transfer all new transaction records to the host.

Possible error codes for the read audit trail command:

ERR_BAD_COMMON_PIN	(Common PIN was incorrect)
ERR_MIAC_NOT_LOCKED	secure module is not locked

Read Group Audit Trail (14H)

Transmit data  
 14H, Group ID, Group PIN  
 Receive data  
 CSB=0 if command successful, appropriate error code otherwise  
 Output length=# of records for group \* 6 if successful, 0 otherwise  
 Output data=audit trail records for group

Notes:

This command is identical to the read audit trail command, except that only records involving the group ID specified in the transmit data are returned to the host. This allows transaction groups to record track their own activities without seeing other groups records.

Possible error codes for the read group audit trail command:

ERR_BAD_GROUP_ID	(Group ID does not exist)
ERR_BAD_GROUP_PIN	(Common PIN was incorrect)
ERR_MIAC_NOT_LOCKED	(The secure module is not locked)

Read Real Time Clock (15H)

Transmit data  
 15H, Common PIN  
 Receive data  
 CSB=0 if the common PIN matches and ERR\_BAD\_COMMON\_PIN otherwise  
 Output length=4  
 Output data=4 most significant bytes of the real time clock

Notes:

This value is not adjusted with a clock offset. This command is normally used by a service provider to compute a clock offset during transaction group creation.

Read Real Time Clock Adjusted (16H)

Transmit data  
 16H, Group ID, Group PIN, ID of offset object  
 Receive data  
 CSB=0 if successful, appropriate error code otherwise  
 Output length=4 if successful, 0 otherwise  
 Output data=Real time clock+clock offset ID

Notes:

This command succeeds if the group ID and group PIN are valid, and the object ID is the ID of a clock offset. The secure module adds the clock offset to the current value of the 4 most significant bytes of the RTC and returns that value in the output data field. Note that a transaction script may be written to perform the same task and put the result in the output data object.

Possible error codes for the real time clock adjusted command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_OBJECT_TYPE	(Object ID is not a clock offset)

Get Random Data (17H)

Transmit data  
 17H, Length (L)  
 Receive data  
 CSB=0 if successful, appropriate error code otherwise  
 Output length=L if successful, 0 otherwise  
 Output data=L bytes of random data if successful

Notes:

This command provides a good source of cryptographically useful random numbers.

Possible error codes for the get random data command are:

ERR\_BAD\_SIZE (Requested number of bytes>128)

Get Firmware Version ID (18H)

Transmit data  
 18H  
 Receive data  
 CSB=0  
 Output length=Length of firmware version ID string  
 Output data=Firmware version ID string

Notes:

This command returns the firmware version ID as a Pascal type string (length+data).

Get Free RAM (19H)

Transmit data  
 19H  
 Receive data  
 CSB=0  
 Output length=2  
 Output data=2 byte value containing the amount of free RAM

Notes:

If the secure module has been locked the output data bytes will both be 0 indicating that all memory not used by transaction groups has been reserved for the audit trail.

Change Group Name (1AH)

Transmit data  
 1AH, Group ID, Group PIN, New Group name  
 Receive data  
 CSB=0 if successful or an appropriate error code otherwise  
 Output length=0  
 Output data=0

Notes:

If the group ID specified exists in the secure module and the PIN supplied is correct, the transaction group name is replaced by the new group name supplied by the host. If a

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group ID of 0 is supplied the PIN transmitted must be the common PIN. If it is correct, the secure module name is replaced by the new name supplied by the host.

Possible error codes for the change group name command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_NAME_LENGTH	(New group name > 16 bytes)

## ERROR CODE DEFINITIONS

## ERR\_BAD\_COMMAND (80H)

This error code occurs when the secure module firmware does not recognize the command just transmitted by the host.

## ERR\_BAD\_COMMON\_PIN (81H)

This error code will be returned when a command requires a common PIN and the PIN supplied does not match the secure module's common PIN. Initially the common PIN is set to 0.

## ERR\_BAD\_GROUP\_PIN (82H)

Transaction groups may have their own PIN, FIG. 6. If this PIN has been set (by a set group PIN command) it must be supplied to access any of the objects within the group. If the Group PIN supplied does not match the actual group PIN, the secure module will return the ERR\_BAD\_GROUP\_PIN error code.

## ERR\_BAD\_PIN\_LENGTH (83H)

There are 2 commands which can change PIN values. The set group PIN and the set common PIN commands. Both of these require the new PIN as well as the old PIN. The ERR\_BAD\_PIN\_LENGTH error code will be returned if the old PIN supplied was correct, but the new PIN was greater than 8 characters in length.

## ERR\_BAD\_OPTION\_BYTE (84H)

The option byte only applies to the common PIN. When the set common PIN command is executed the last byte the host supplies is the option byte (described in command section). If this byte is unrecognizable to the secure module, it will return the ERR\_BAD\_OPTION\_BYTE error code.

## ERR\_BAD\_NAME\_LENGTH (85H)

When the create transaction group command is executed, one of the data structures supplied by the host is the group's name. The group name may not exceed 16 characters in length. If the name supplied is longer than 16 characters, the ERR\_BAD\_NAME\_LENGTH error code is returned.

## ERR\_INSUFFICIENT\_RAM (86H)

The create transaction group and create object commands return this error code when there is not enough heap available in the secure module.

## ERR\_MIAC\_LOCKED (87H)

When the secure module has been locked, no groups or objects can be created or destroyed. Any attempts to create or delete objects will generate an ERR\_MIAC\_LOCKED error code.

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## ERR\_MIAC\_NOT\_LOCKED (88H)

If the secure module has not been locked there is no audit trail. If one of the audit trail commands is executed this error code will be returned.

## ERR\_GROUP\_LOCKED (89H)

Once a transaction group has been locked object creation within that group is not possible. Also the objects attributes and types are frozen. Any attempt to create objects or modify their attribute or type bytes will generate an ERR\_GROUP\_LOCKED error code.

## ERR\_BAD\_OBJECT\_TYPE (8AH)

When the host sends a create object command to the secure module, one of the parameters it supplies is an object type (see command section). If the object type is not recognized by the firmware it will return an ERR\_BAD\_OBJECT\_TYPE error code.

## ERR\_BAD\_OBJECT\_ATTR (8BH)

When the host sends a create object command to the secure module, one of the parameters it supplies is an object attribute byte (see command section). If the object attribute byte is not recognized by the firmware it will return an ERR\_BAD\_OBJECT\_ATTR error code.

## ERR\_BAD\_SIZE (8CH)

An ERR\_BAD\_SIZE error code is normally generated when creating or writing an object. It will only occur when the object data supplied by the host has an invalid length.

## ERR\_BAD\_GROUP\_ID (8DH)

All commands that operate at the transaction group level require the group ID to be supplied in the command packet. If the group ID specified does not exist in the secure module it will generate an ERR\_BAD\_GROUP\_ID error code.

## ERR\_BAD\_OBJECT\_ID (8EH)

All commands that operate at the object level require the object ID to be supplied in the command packet. If the object ID specified does not exist within the specific transaction group (also specified in the command packet) the secure module will generate an ERR\_BAD\_OBJECT\_ID error code.

## ERR\_INSUFFICIENT\_FUNDS (8FH)

If a script object that executes financial transactions is invoked and the value of the money register is less than the withdrawal amount requested an ERR\_INSUFFICIENT\_FUNDS error code will be returned.

## ERR\_OBJECT\_LOCKED (90H)

Locked objects are read only. If a write object command is attempted and it specifies the object ID of a locked object the secure module will return an ERR\_OBJECT\_LOCKED error code.

## ERR\_OBJECT\_PRIVATE (91H)

Private objects are not directly readable or writable. If a read object command or a write object command is attempted, and it specifies the object ID of a private object, the secure module will return an ERR\_OBJECT\_PRIVATE error code.

## ERR\_OBJECT\_DESTROYED (92H)

If an object is destructible and the transaction group's destructor is active the object may not be used by a script. If a script is invoked which uses an object which has been destructed, an ERR\_OBJECT\_DESTROYED error code will be returned by the secure module.

The exemplary embodiment of the present invention is preferably placed within a durable stainless steel, token-like can. It is understood that an exemplary secure module can be placed in virtually any articulatable item. Examples of articulatable items include credit cards, rings, watches, wallets, purses, necklaces, jewelry, ID badges, pens, clipboards, etc.

The secure module 108 preferably is a single chip "trusted computer". By the word "trusted" it is meant that the computer is extremely secure from tampering by unwarranted means. The secure module incorporates a numeric coprocessor optimized for math intensive encryption. The BIOS is preferably immune to alteration and specifically designed for very secure transactions.

Each secure module can have a random "seed" generator with the ability to create a private/public key set. The private key never leaves the secure module and is only known by the secure module. Furthermore, discovery of the private key is prevented by active self-destruction upon wrongful entry into the secure module. The secure module can be bound to the user by a personal identification number (PIN).

When transactions are performed by the secure module 108 certificates of authentication are created by either or both the secure module and a system the secure module communicates with. The certificate can contain a variety of information. In particular, the certificate may contain:

- 1) who is the secure module user via a unique registration number and a certified public key.
- 2) when the transaction took place via a true-time stamping of the transaction.
- 3) where the transaction took place via a registered secure module interface site identification.
- 4) security information via uniquely serialized transactions and digital sign on message digests.
- 5) secure module status indicated as valid, lost, or expired.

Although a preferred embodiment of the method and apparatus of the present invention has been illustrated in the accompanying Drawings and described in the foregoing Detailed Description, it will be understood that the invention is not limited to the embodiment disclosed, but is capable of numerous rearrangements, modifications and substitutions without departing from the spirit of the invention as set forth and defined by the following claims.

What is claimed is:

1. A system for communicating data securely, comprising:
  - a first portable module comprising:
    - a nonvolatile memory for storing a first data;
    - a first real time clock circuit for time stamping data transactions;
    - a counter for counting a transaction count;
    - an input/output circuit;
    - a substantially unique electronically readable identification number readable by said input/output circuit; and
    - a memory control circuit in electrical communication with said nonvolatile memory, said real time clock, said counter, and said input/output circuit;
  - a portable module reader that can be placed in communication with said first portable module, said portable module reader can be connected to a plurality of other devices;
  - a secure microcontroller based module in electronic communication with said portable module reader, said secure microcontroller comprising:
    - a microcontroller core;
    - a math coprocessor, in communication with said microcontroller core, for processing encryption calculations;
    - an energy circuit for storing energy;
    - a memory circuit connected to said microcontroller core;
    - a memory circuit in communication with said microcontroller core; and
    - a second real time clock circuit in communication with said microcontroller,
 said combination of said portable module reader and said secure microcontroller performing secure data transfers with said first portable module.
2. The system for communicating data securely of claim 1, wherein said plurality of other devices includes at least one of a credit card reader, a cash machine, an automatic teller machine, and a phone line.
3. The system for communicating data securely of claim 1, wherein said first data is a packet of encrypted data.
4. The system for communicating data securely of claim 1, wherein said first portable module communicates with said portable module reader via a single wire bus comprising a single bidirectional communication wire and a ground connection.
5. The system for communicating data securely of claim 1, wherein said first module can create random public/private key sets for encryption purposes.
6. The system for communicating data securely of claim 1, wherein said secure microcontroller can create random public/private key sets for encryption purposes.

\* \* \* \* \*

UNITED STATES PATENT AND TRADEMARK OFFICE  
**CERTIFICATE OF CORRECTION**

PATENT NO. : 5,940,510  
DATED : Aug. 17, 1999  
INVENTOR(S) : Curry et al.

It is certified that error appears in the above-identified patent and that said Letters Patent is hereby corrected as shown below:

On the title page, item:

Item [57], line 1            Replace "rotates"  
   With --relates--

Signed and Sealed this

Twenty-second Day of February, 2000

Attest:



Q. TODD DICKINSON

Attesting Officer

Commissioner of Patents and Trademarks

# **EXHIBIT D**



US005949880A

**United States Patent** [19]  
**Curry et al.**

[11] **Patent Number:** **5,949,880**  
[45] **Date of Patent:** **Sep. 7, 1999**

[54] **TRANSFER OF VALUABLE INFORMATION BETWEEN A SECURE MODULE AND ANOTHER MODULE**

[75] Inventors: **Stephen M. Curry**, Dallas; **Donald W. Loomis**, Coppell; **Michael L. Bolan**, Dallas, all of Tex.

[73] Assignee: **Dallas Semiconductor Corporation**, Dallas, Tex.

[21] Appl. No.: **08/978,798**

[22] Filed: **Nov. 26, 1997**

**Related U.S. Application Data**

[62] Division of application No. 08/594,975, Jan. 31, 1996.

[51] **Int. Cl.<sup>6</sup>** ..... **H04L 9/00**

[52] **U.S. Cl.** ..... **380/24; 380/25; 705/39; 705/42**

[58] **Field of Search** ..... **380/23, 24, 25; 705/39, 40, 42**

[56] **References Cited**

**U.S. PATENT DOCUMENTS**

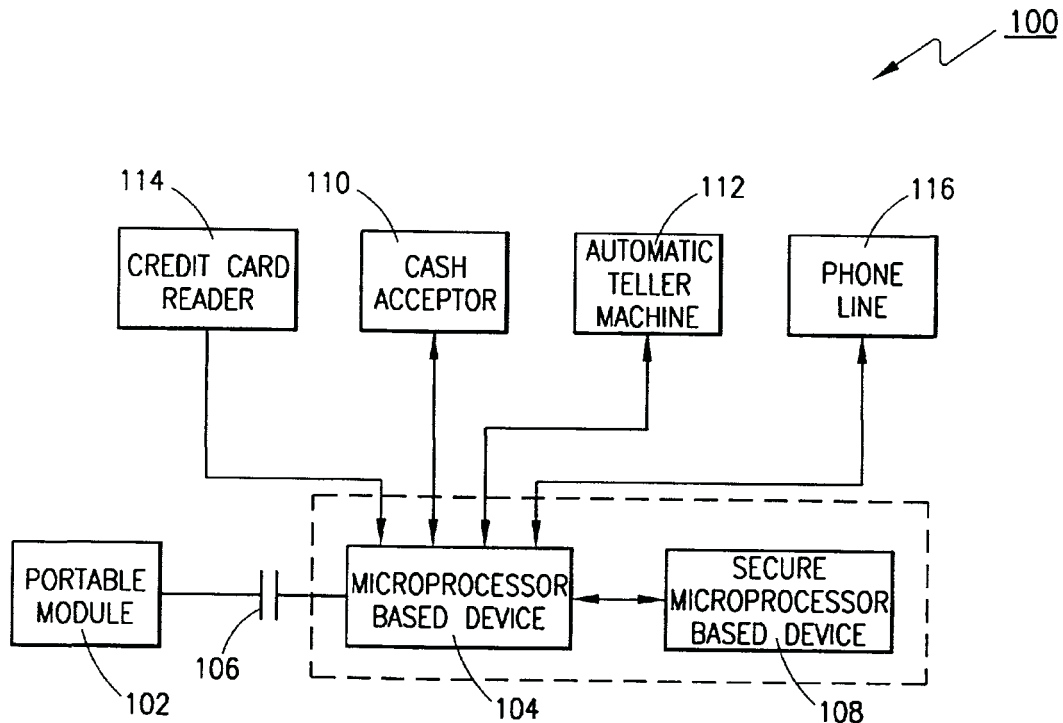
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5,642,419	6/1997	Rosen	380/23
5,671,280	9/1997	Rosen	380/24

*Primary Examiner*—Thomas H. Tarza  
*Assistant Examiner*—Carmen D. White  
*Attorney, Agent, or Firm*—Jenkins & Gilchrist

[57] **ABSTRACT**

The present invention relates to system, apparatus and method for communicating valuable data from a portable module to another module via an electronic device. More specifically, the disclosed system, apparatus and method are useful for enabling a user to fill a portable module with a cash equivalent and to spend the cash equivalent at a variety of locations. The disclosed system incorporates an encryption/decryption method.

**6 Claims, 8 Drawing Sheets**



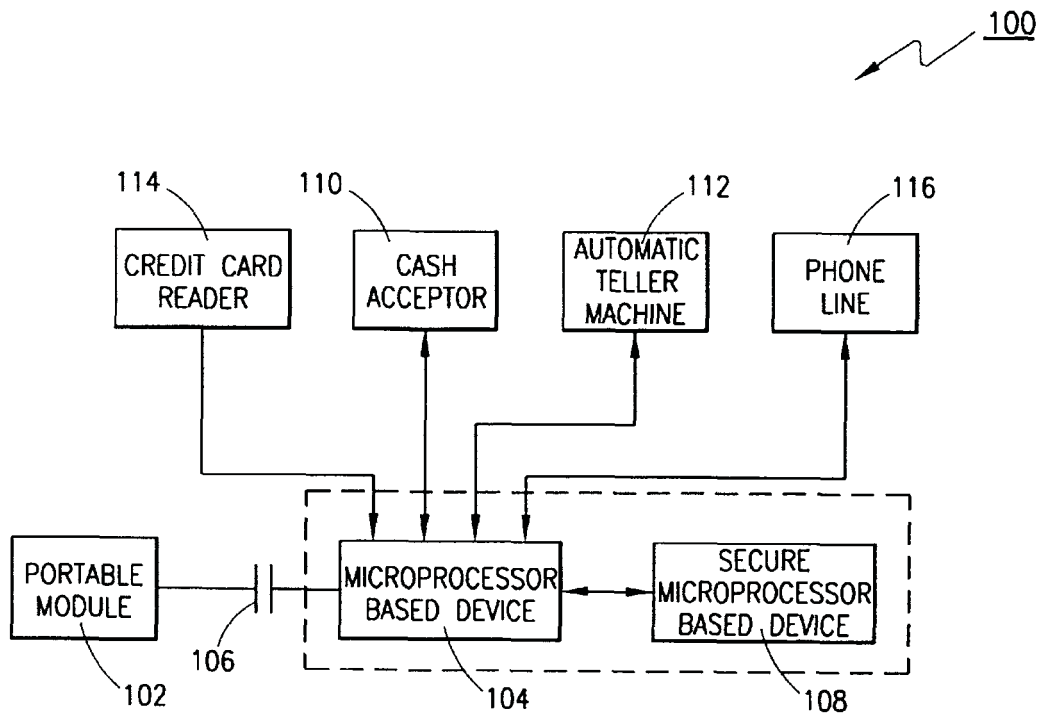


FIG. 1



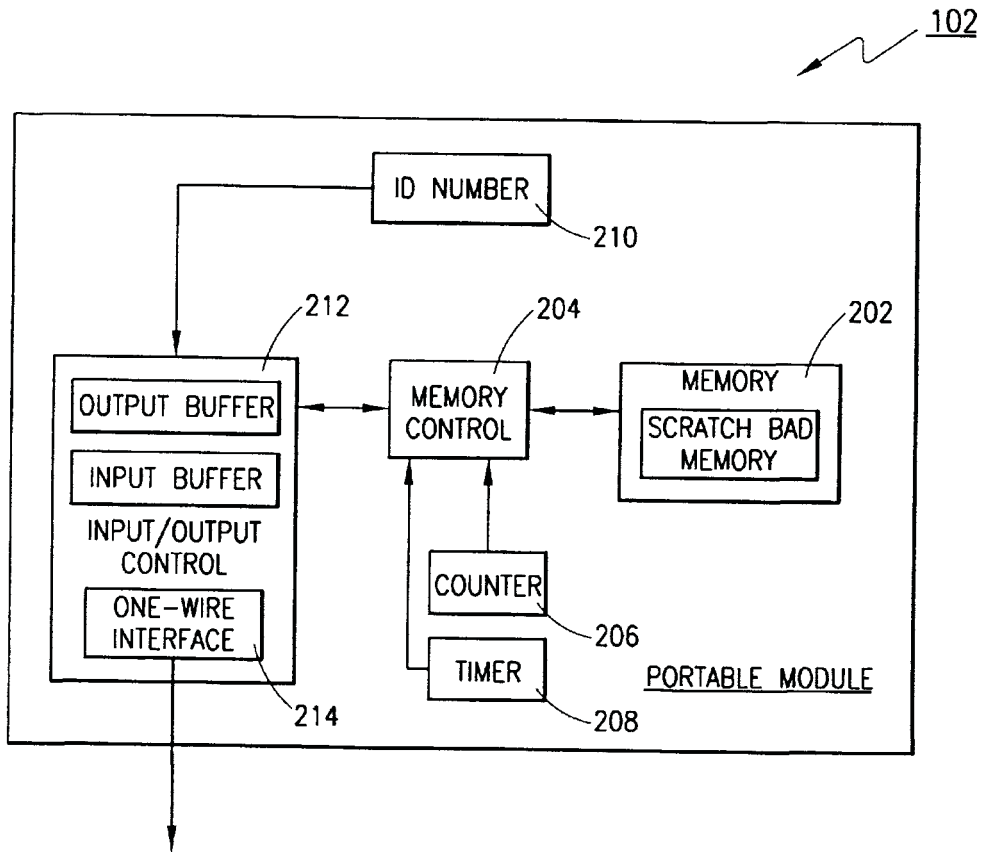


FIG. 2

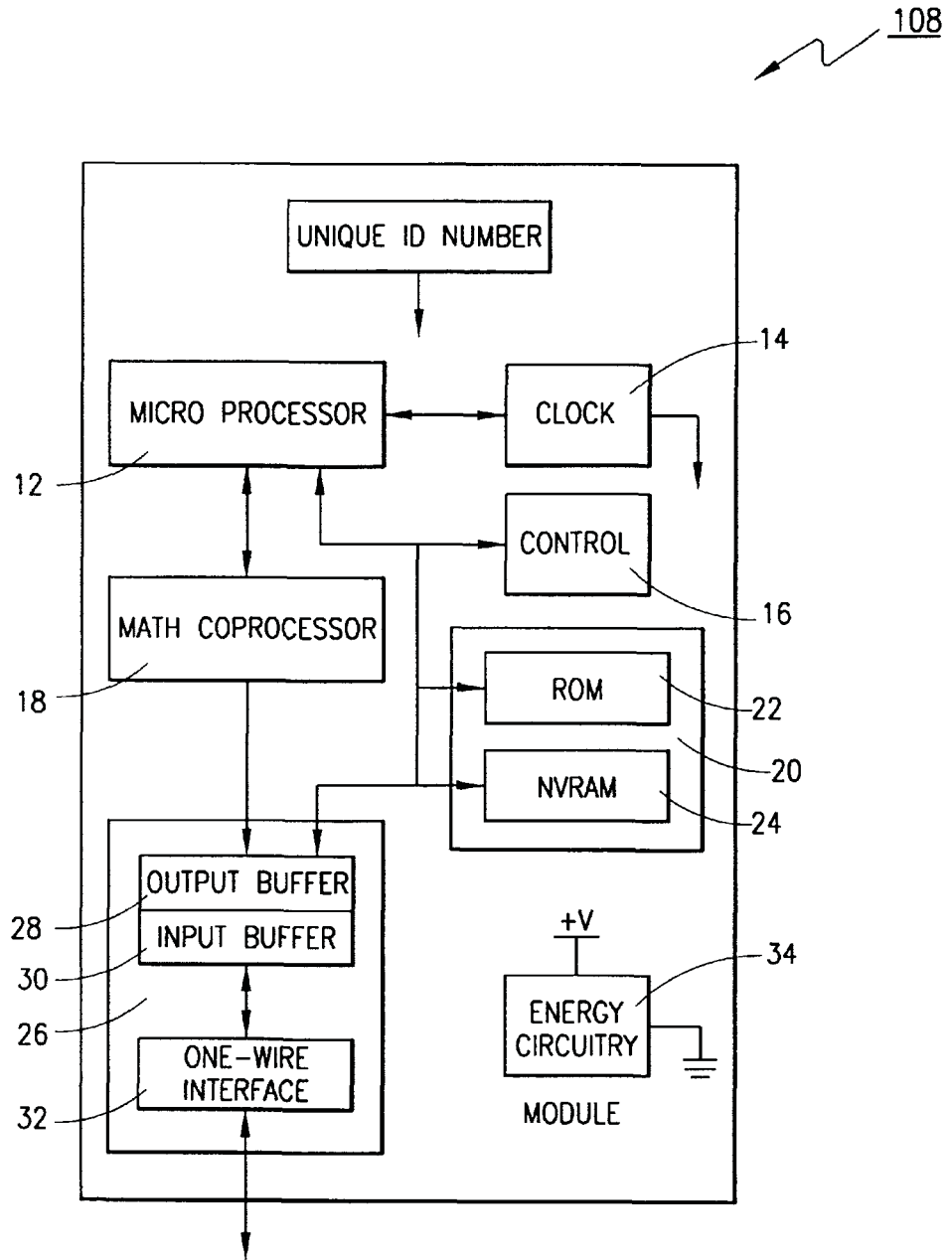


FIG. 3

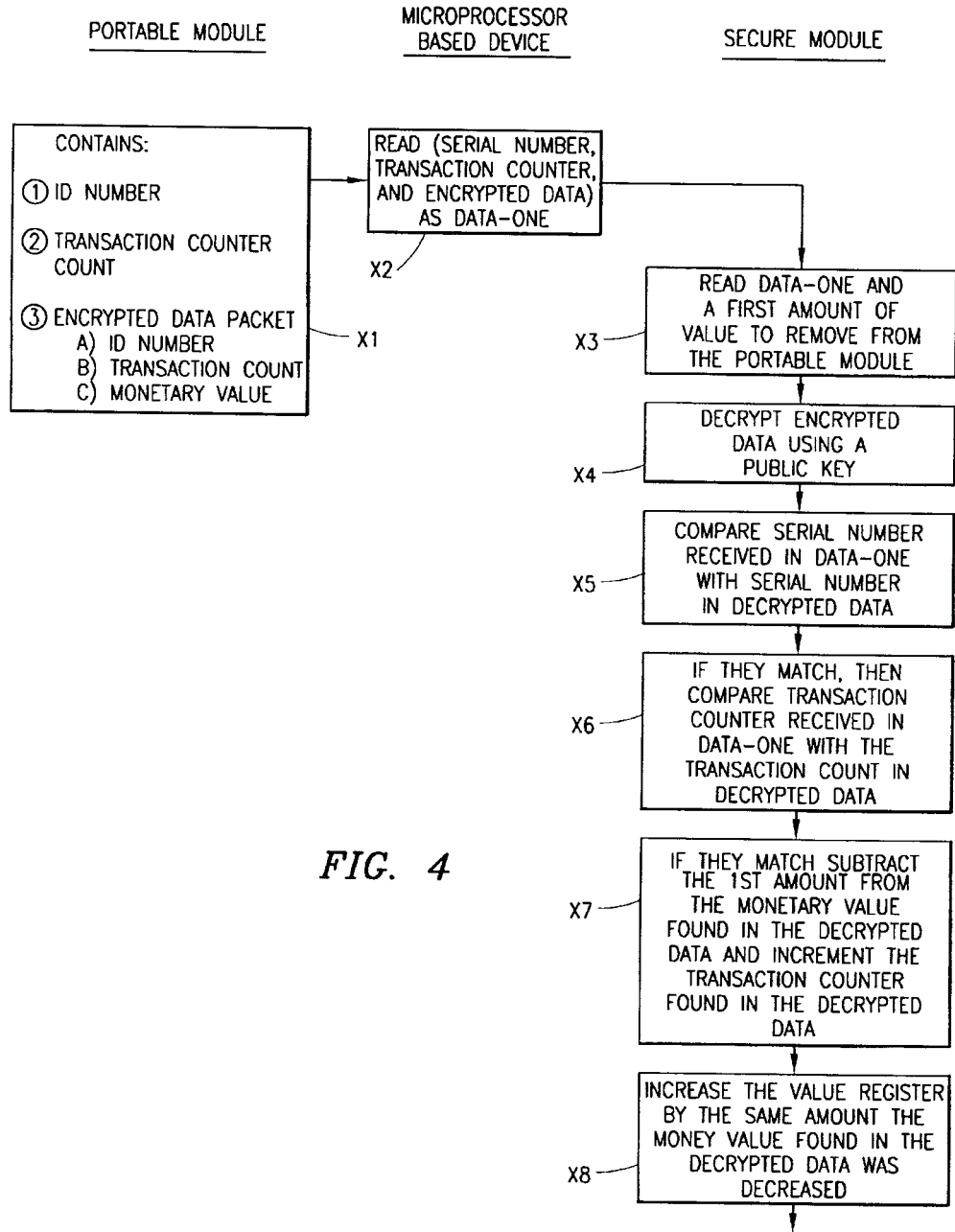
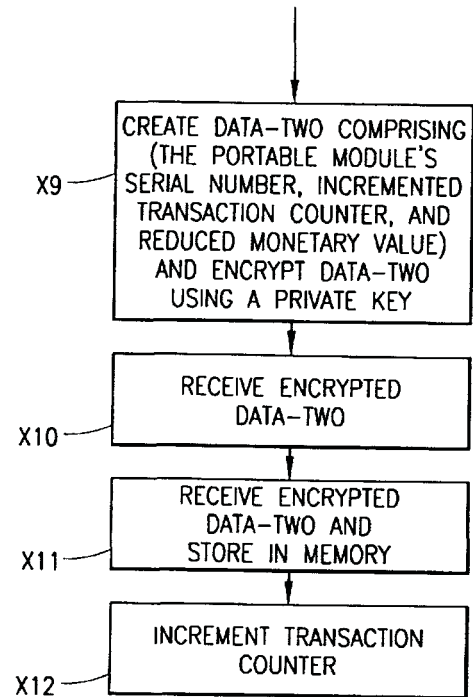


FIG. 4

PORTABLE MODULE

MICROPROCESSOR  
BASED DEVICE

SECURE MODULE



**FIG. 4**  
(CONTINUED)

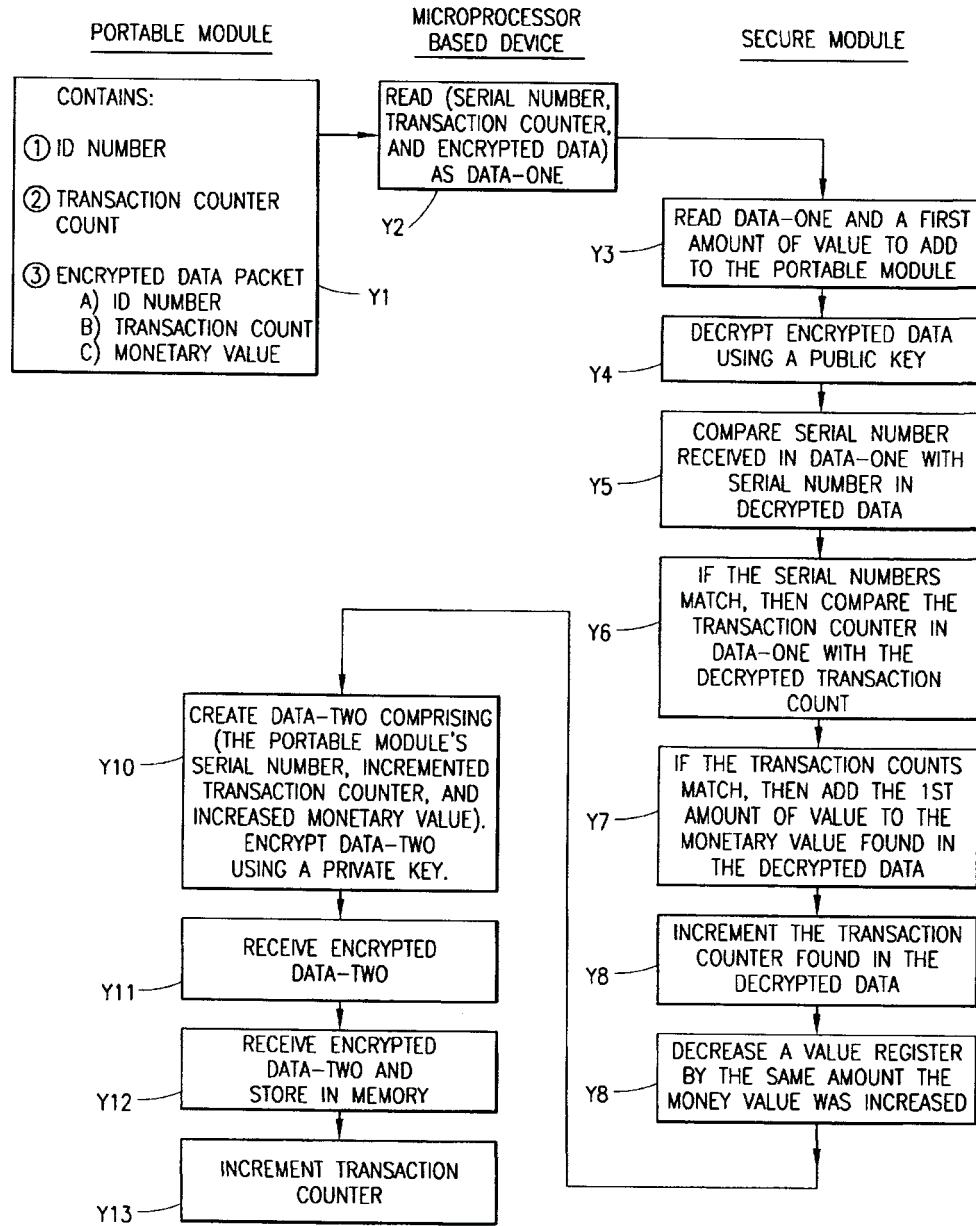


FIG. 5

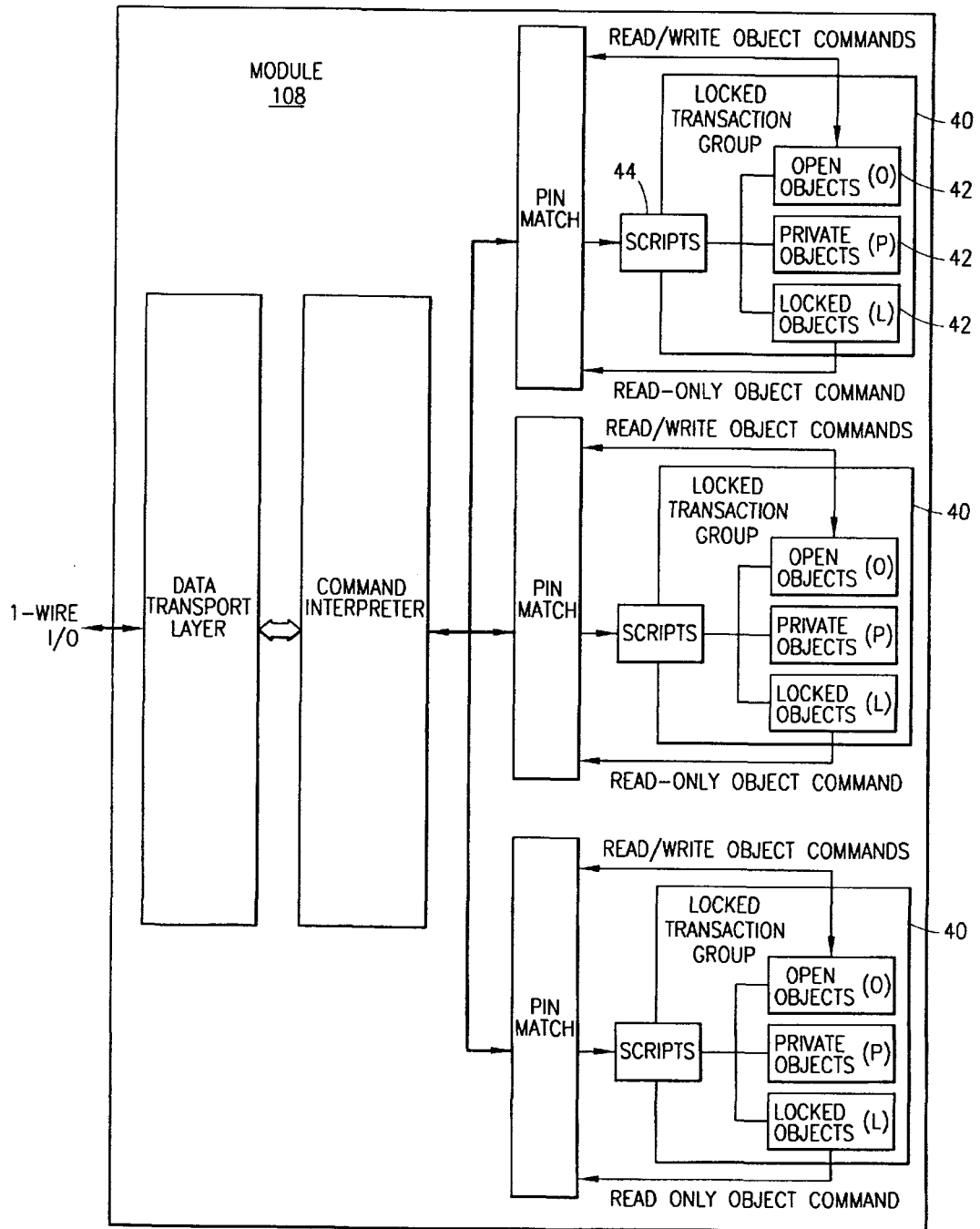


FIG. 6

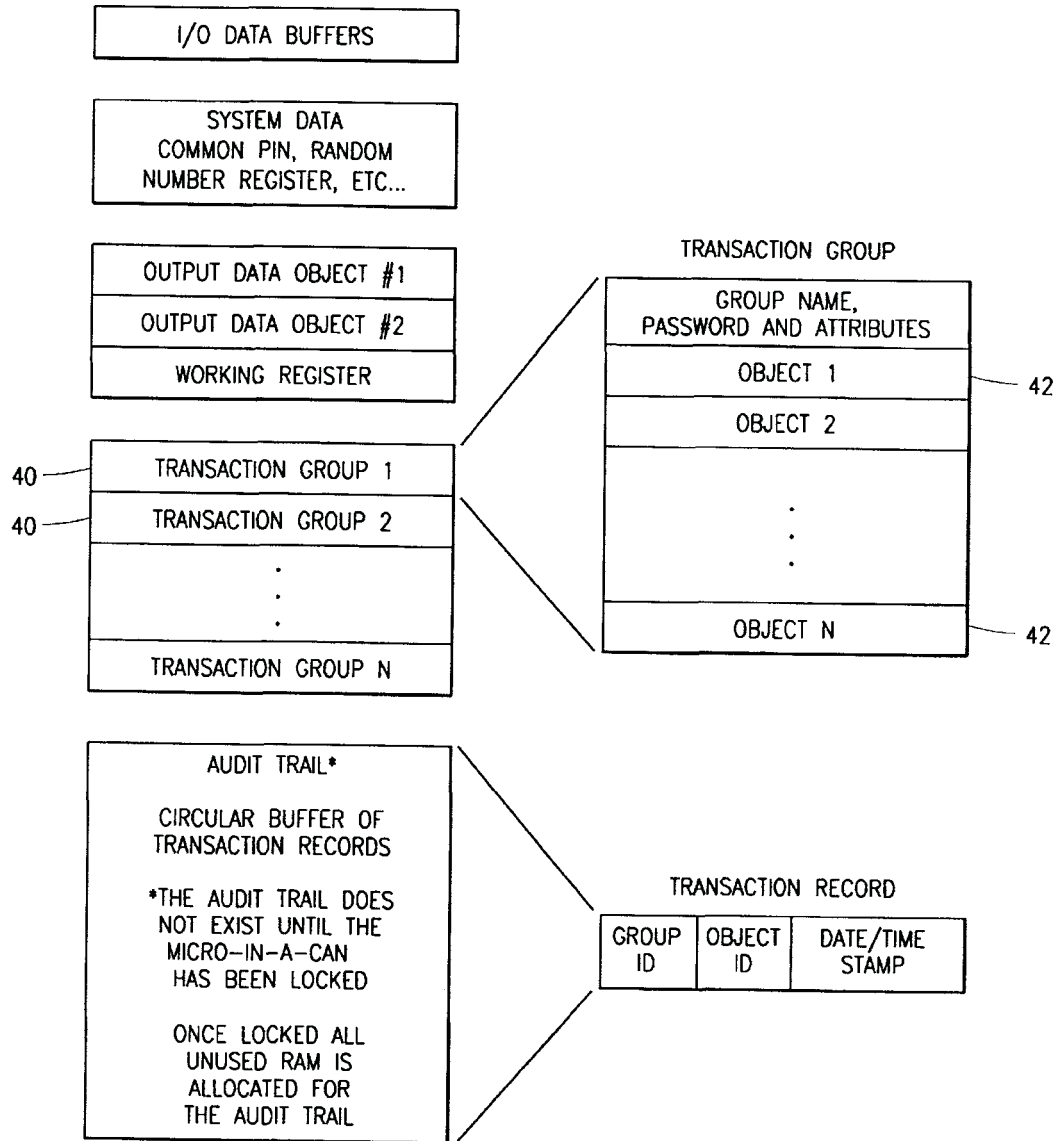


FIG. 7

## TRANSFER OF VALUABLE INFORMATION BETWEEN A SECURE MODULE AND ANOTHER MODULE

This application is a Divisional of application Ser. No. 08/594,975 filed on Jan. 31, 1996.

### CROSS REFERENCE TO OTHER APPLICATIONS

The following applications of common assignee contains related subject matter and is hereby incorporated by reference:

Ser. No. UNKNOWN, filed Jan. 31, 1996, entitled METHOD, APPARATUS, SYSTEM AND FIRMWARE FOR SECURE TRANSACTIONS; and

Ser. No. UNKNOWN, filed Jan. 31, 1996, entitled METHOD, APPARATUS AND SYSTEM FOR TRANSFERRING UNITS OF VALUE.

### BACKGROUND OF THE INVENTION

#### 1. Technical Field of the Invention

The present invention relates to a method and system for transferring valuable information securely between a secure module and another module. More particularly, the present invention relates to transferring units of value between a microprocessor based secure module and another module used for carrying a monetary equivalent.

#### 2. Description of Related Art

In the past the preferred means for paying for an item was cash. As our society has become more advanced, credit cards have become an accepted way to pay for merchandise or services. The payment is not a payment to the merchant, but instead is a credit given by a bank to the user that the merchant accepts as payment. The merchant collects money from the bank based on the credit. As time goes on, cash is used less and less, and money transfers between parties are becoming purely electronic.

Present credit cards have magnetic strips to identify the owner of the card and the credit provider. Some credit cards have electronic circuitry installed that identifies the credit card owner and the credit or service provider (the bank).

The magnetic strips installed in present credit cards do not enable the card to be used as cash. That is the modern credit card does not allow the consumer to buy something with the credit card and the merchant to receive cash at the time of the transaction. Instead, when the consumer buys something on credit, the merchant must later request that the bank pay for the item that the consumer bought. The bank then bills the consumer for the item that was bought.

Thus, there is a need for an electronic system that allows a consumer to fill an electronic module with a cash equivalent in the same way a consumer fills his wallet with cash. When the consumer buys a product or service from a merchant, the consumer's module can be debited and the merchant's cash drawer can be credited without any further transactions with a bank or service provider.

### SUMMARY OF THE INVENTION

The present invention is an apparatus, system and method for communicating a cash equivalent electronically to and from a portable module. The portable module can be used as a cash equivalent when buying products and services in the market place.

The present invention comprises a portable module that can communicate to a secure module via a microprocessor

based device. The portable module can be carried by a consumer, filled with electronic money at an add-money station, and be debited by a merchant when a product or service is purchased by the consumer. As a result of a purchase, the merchant's cash drawer will indicate an increase in cash value.

### BRIEF DESCRIPTION OF THE DRAWINGS

A more complete understanding of the method and apparatus of the present invention may be had by reference to the following Detailed Description when taken in conjunction with the accompanying Drawings wherein:

FIG. 1 depicts an exemplary system for transferring valuable information between a module and a secure device;

FIG. 2 is a block diagram of an embodiment of a portable module;

FIG. 3 is a block diagram of an embodiment of a microprocessor based module;

FIG. 4 is an exemplary technique for transferring valuable data securely into a portable module;

FIG. 5 is an exemplary technique for transferring valuable data securely out of a portable module;

FIG. 6 is an exemplary organization of the software and firmware within a secure microprocessor based device; and

FIG. 7 is an exemplary configuration of software and firmware within a secure microprocessor based device.

### DETAILED DESCRIPTION OF A PRESENTLY PREFERRED EXEMPLARY EMBODIMENT

FIG. 1 depicts a block diagram of an exemplary system **100** for transferring valuable information to and from a portable module. A portable module **102**, which will be described in more detail later, communicates to a microprocessor based device **104**. The portable module **102** may contain information that represents units of exchange or a currency equivalent. The microprocessor based device **104** can be any of an unlimited number of devices. For example, the microprocessor based device **104** could be a personal computer, an add-a-fare machine at a train or bus station (similar to those in today's District of Columbia metro stations), a turn style, a toll booth, a bank's terminal, a ride at a carnival, a washing machine at a Laundromat, a locking device, a mail metering device or any device that controls access, or meters a monetary equivalent, etc.

The means for communication **106** between the portable module **102** and the microprocessor based device **104** is preferably via a single wire or contact connection. The single wire connection **106** preferably incorporates a communication protocol that allows the portable module **102** and the microprocessor based device **104** to communicate in a bidirectional manner. Preferably the communication protocol is a one-wire protocol developed by Dallas Semiconductor. It is understood that the means for communicating **106** is not limited to a single wire connection. The communication means **106** could be multiple wires, a wireless communication system, infrared light, any electromagnetic means, a magnetic technique, or any other similar technique.

The microprocessor based device **104** is electrically connected to another microprocessor based device, which is preferably a secure device **108**. The term secure device means that the device is designed to contain a secret code and the secret code is extremely difficult to learn. An example of a secure device **108** is explained later in this document.

The microprocessor based device **104** can be connected to a variety of other devices. Such devices include, but are not



limited to a cash acceptor **110**, an automatic teller machine (ATM) **112**, a credit card reader **114**, and a phone line **116**.

The cash acceptor **110** is adapted to receive cash in the form of currency, such as dollar bills or coins. The cash acceptor **110**, preferably, determines the value of the accepted currency. The cash acceptor **110** communicates to the microprocessor based device **104** and informs the device **104** of how much currency has been deposited in the cash acceptor **110**.

The cash acceptor **110** can also be a device which provides currency. That is, the cash acceptor **110** in response to a communication from the microprocessor based device **104**, may provide a metered amount of currency to a person.

The credit card reader **114**, and ATM **112** can also be attached to the microprocessor based device **104**. The credit card reader **114** could be used to read a user's credit card and then, when authorized, either communicate to the microprocessor based device **104** that units of exchange need to be added to the portable module or that units of exchange need to be extracted from the portable module to pay for a good, service or credit card bill.

The ATM **112** may also be connected to the microprocessor based device. Via communications from the ATM **112**, the microprocessor based device **104** can be informed that units of exchange need to be added or subtracted from the portable module **102**.

Furthermore, it is also possible that the microprocessor based device **104** is connected to a phone line **116**. The phone line may be used for a variety of things. Most importantly, the phone line may be used to allow the microprocessor based device **104** to communicate with a network of devices. Such telephonic communication may be for validating transactions or for aiding the accounting of transactions that are performed via the microprocessor based device's **104** aid. It is further understood that the phone line may be any of a vast variety of communication lines including wireless lines. Video, analog, or digital information may be communicated over the phone line **116**.

FIG. 2 depicts a preferred exemplary portable module **102**. The portable module **102** is preferably a rugged read/write data carrier that can act as a localized data base and be easily accessed with minimal hardware. The module can be incorporated in a vast variety of portable items which includes, but is not limited to a durable micro-can package that is highly resistant to environmental hazards such as dirt, moisture, and shock. The module can be incorporated into any object that can be articulated by a human or thing, such as a ring, bracelet, wallet, name tag, necklace, baggage, machine, robotic device, etc. Furthermore, the module **102** could be attached to a stationary item and the microprocessor based device **104** may be articulated to the portable module **102**. For example, the module **102** may be attached to a piece of cargo and a module reader may be touched to or brought near the module **102**. The module reader may be part of the microprocessor based device **104**.

The portable module **102** comprises a memory **202** that is preferably, at least in part, nonvolatile memory for storing and retrieving vital information pertaining to the system to which the module **102** may become attached to. The memory **202** may contain a scratchpad memory which may act as a buffer when writing into memory. Data is first written to the scratchpad where it can be read back. After data has been verified, the data is transferred into the memory.

The module **102** also comprises a counter **206** for keeping track of the number of transactions the module has per-

formed (the number of times certain data in the memory of the module has been changed). A timer **102** may be provided in the module to provide the ability to time stamp transactions performed by the module. A memory controller **204** controls the reading and writing of data into and out of the memory **202**.

The module also may comprise an identification number **210**. The identification number preferably uniquely identifies the portable module from any other portable module.

An input/output control circuit **212** controls the data flow into and out of the portable module **102**. The input/output control ("I/O") **212** preferably has an input buffer and an output buffer and interface circuitry **214**. As stated above, the interface circuitry **214** is preferably a one-wire interface. Again, it is understood that a variety of technologies can be used to interface the portable module **102** to another electronic device. A single wire or single connection is preferred because the mechanics of making a complete connection is simplified. It is envisioned that a proximity/wireless communication technique is also a technique for communicating between the module **102** and another device. Thus, the interface circuit **214** can be a single wire, multiple wire, wireless, electromagnetic, magnetic, light, or proximity, interface circuit.

FIG. 3 depicts a block diagram of an exemplary secure microprocessor based device ("secure device") **108**. The secure device circuitry can be a single integrated circuit. It is understood that the secure device **108** could also be a monolithic or multiple circuits combined together. The secure device **108** preferably comprises a microprocessor **12**, a real time clock **14**, control circuitry **16**, a math coprocessor **18**, memory circuitry **20**, input/output circuitry **26**, and an energy circuit **34**.

The secure device **108** could be made small enough to be incorporated into a variety of objects including, but not limited to a token, a card, a ring, a computer, a wallet, a key fob, a badge, jewelry, a stamp, or practically any object that can be grasped and/or articulated by a user of the object. In the present system **100**, the secure device **108** is preferably adapted to be a trusted certifying authority. That is the secure device **108** is a trusted computer. The secure device **108** comprises a numeric coprocessor **18** optimized for math intensive encryption. The BIOS is immune to alteration and is specifically designed for secure transactions. This secure device **108** is preferably encased in a durable, dirt, moisture and shock resistant stainless steel enclosure, but could be encased in wide variety of structures so long as specific contents of the secure device **108** are extremely difficult to decipher. The secure device **108**. The secure device **108** may have the ability to store or create a private/public key set, whereby the private key never leaves the secure device **108** and is not revealed under almost any circumstance. Furthermore, the secure module **108** is preferably designed to prevent discovery of the private key by an active self-destruction of the key upon wrongful entry.

The microprocessor **12** is preferably an 8-bit microprocessor, but could be 16, 32, 64 or any operable number of bits. The clock **14** provides timing for the module circuitry. There can also be separate clock circuitry **14** that provides a continuously running real time clock.

The math coprocessor circuitry **18** is designed and used to handle very large numbers. In particular, the coprocessor will handle the complex mathematics of RSA encryption and decryption or other types of math intensive encryption or decryption techniques.

The memory circuitry **20** may contain both read-only-memory and non-volatile random-access-memory.

Furthermore, one of ordinary skill in the art would understand that volatile memory, EPROM, SRAM and a variety of other types of memory circuitry might be used to create an equivalent device.

Control circuitry 16 provides timing, latching and various necessary control functions for the entire circuit.

An input/output circuit 26 enables bidirectional communication with the secure module 108. The input/output circuitry 26 preferably comprises at least an output buffer and an input buffer. For communication via a one-wire bus, one-wire interface circuitry can be included with the input/output circuitry 26. It is understood that the input/output circuitry 26 of the secure device 108 can be designed to operate on a single wire, a plurality of wires or any means for communicating information between the secure module 108 and the microprocessor based device 104.

An energy circuit 34 may be necessary to maintain stored information in the memory circuitry 20 and/or aid in powering the other circuitry in the module 108. The energy circuit 34 could consist of a battery, capacitor, R/C circuit, photo-voltaic cell, or any other equivalent energy producing circuit or means.

The firmware architecture of the secure module 108 and how it operates within the exemplary system for transferring valuable information, such as units of exchange or currency, between the secure module 108 and a portable module 102 will now be discussed. The secure module 108 provides encryption and decryption services for confidential data transfer through the microprocessor based device 104. The following examples are intended to illustrate a preferred feature set of the secure module 108 and to explain the services that the exemplary system 100 can offer. These applications and examples by no means limit the capabilities of the invention, but instead bring to light a sampling of its capabilities.

#### I. Overview of the Preferred Secure Module 108 and its Firmware Design

Referring to FIG. 3 again, the secure module 108 preferably contains a general-purpose, 8051-compatible micro controller 12 or a reasonably similar product, a continuously running real-time clock 14, a high-speed modular exponentiation accelerator for large integers (math coprocessor) 18, input and output buffers 28, 30 with a one-wire interface 32 for sending and receiving data, 32 Kbytes of ROM memory 22 with preprogrammed firmware, 8 Kbytes of NVRAM (non-volatile RAM) 24 for storage of critical data, and control circuitry 16 that enables the micro controller 12 to be powered up to interpret and act on the data placed in an input data object. The module 108 draws its operating power from a single wire, one-wire communication line. The micro controller 12, clock 14, memory 20, buffers 28, 30, one-wire front-end 32, modular exponentiation accelerator 18, and control circuitry 16 are preferably integrated on a single silicon chip and packaged in a stainless steel micro can using packaging techniques which make it virtually impossible to probe the data in the NVRAM 24 without destroying the data. Initially, most of the NVRAM 24 is available for use to support applications such as those described below. One of ordinary skill will understand that there are many comparable variations of the module design. For example, volatile memory might be used, or an interface other than a one-wire interface could be used.

The secure module 108 is preferably intended to be used first by a Service Provider who loads the secure module 108 with data to enable it to perform useful functions, and second by an End User who issues commands to the secure module 108 to perform operations on behalf of the Service

Provider for the benefit of the End User. For this reason, the secure module 108 offers functions to support the Service Provider in setting up the module for an intended application. It also offers functions to allow the End User to invoke the services offered by the Service Provider.

Each Service Provider can reserve a block of NVRAM memory to support its services by creating a transaction group 40 (refer to FIGS. 6 and 7). A transaction group 40 is simply a set of software objects 42 that are defined by the Service Provider. These objects 42 include both data objects (encryption keys, transaction counts, money amounts, date/time stamps, etc.) and transaction scripts 44 which specify how to combine the data objects in useful ways. Each Service Provider creates his own transaction group 40, which is independent of every other transaction group 40. Hence, multiple Service Providers can offer different services in the same module 108. The number of independent Service Providers that can be supported depends on the number and complexity of the objects 42 defined in each transaction group 40. Examples of some of the objects 42 that can be defined within a transaction group 40 are the following:

RSA Modulus	Clock Offset
RSA Exponent	Random SALT
Transaction Script	Configuration Data
Transaction Counter	Input Data
Money Register	Output Data
Destructor	

Within each transaction group 40 the secure module 108 will initially accept certain commands which have an irreversible effect. Once any of these irreversible commands are executed in a transaction group 40, they remain in effect until the end of the module's useful life or until the transaction group 40, to which it applies, is deleted from the secure module 108. In addition, there are certain commands which have an irreversible effect until the end of the module's life or until a master erase command is issued to erase the entire contents of the secure module 108. These commands will be discussed further below. These commands are essential to give the Service Provider the necessary control over the operations that can be performed by the End User. Examples of some of the irreversible commands are:

Privatize Object	Lock Object
Lock Transaction Group	Lock Micro-In-A-Can™

Since much of the module's utility centers on its ability to keep a secret, the Privatize command is a very important irreversible command.

Once the secure module 108, as a whole, is locked, the remaining NVRAM memory 24 is allocated for a circular buffer for holding an audit trail of previous transactions. Each of the transactions are identified by the number of the transaction group, the number of objects 42 within the specified group, and the date/time stamp.

The fundamental concept implemented by the firmware is that the Service Provider can store transaction scripts 44 in a transaction group 40 to perform only those operations among objects that he wishes the End User to be able to perform. The Service Provider can also store and privatize RSA key or keys (encryption keys) that allow the secure module 108 to "sign" transactions on behalf of the Service Provider, thereby guaranteeing their authenticity. By privatizing and/or locking one or more objects 42 in the trans-

action group 40, the Service Provider maintains control over what the secure module 108 is allowed to do on his behalf. The End User cannot add new transaction scripts 44 and is therefore limited to the operations on objects 42 that can be performed with the transaction scripts 44 programmed by the Service Provider.

## II. Usage Models of the Secure Module 108 and Portable Module 102

This section presents practical applications of the system 100. Each of these applications is described in enough detail to make it clear why the secure module 108 and portable module 102 are important to the system application.

### A. Transferring Units of Exchange Out of a Portable Module 102

This section describes an example of how a portable module 102 and a secure module 108 operate in conjunction with the microprocessor based device 104 so that units of exchange can be securely transferred out of the portable module 102 and deposited into the secure module 108 and/or potentially communicated to at least one of the cash acceptor 110, ATM 112, credit card reader 114, or the phone line 116.

Referring to FIG. 4, initially the portable module 102 contains its ID number, a count within its transaction counter and an encrypted data packet stored in memory. Encrypted within the data packet is the portable modules ID number, the portable modules transaction count number, and the amount of value (the monetary value) of the portable module at the present time X1.

The user of the portable module touches, or somehow puts the portable module 102 into communication with the microprocessor based device 104. For explanation purposes, suppose the portable module 102 is being used as a token used to pay for a train fare. Thus, the microprocessor based device 104 could be, in this case, a turn style that allows the user to enter a train platform. The cost of entering the train platform is known by the microprocessor based device 104.

The microprocessor based device 104 reads the portable module's serial number, transaction count, and the encrypted data packet X2. This data could be referred to as a first data.

The microprocessor device 104 then provides the first data along with a first value, being the amount of value to be debited from the portable token (the train fare), to the secure module 108 X3. The secure module 108 decrypts the encrypted data found in the first data using a public key X4.

Next, the secure module 108 makes a few comparisons to make sure that the data received is good data and not counterfeit. The secure module 108 compares the serial number received in the first data with the decrypted serial number X5. If the two serial numbers match then the secure module 108 compares the transaction count received in the first data with the decrypted transaction count X6. If the two transaction counts match then the secure module is comfortable that the data received is not counterfeit data. It is understood that the comparisons can be done in any order.

Furthermore, there may have been a time stamp sent from the portable module 102. The time stamp may indicate a variety of things. One thing could be an indication of whether the portable module is still valid or the time stamp may further enable the secure module to decide if the data is or is not counterfeit.

Assuming all the data passed to the secure module 108 is determined to be valid data, the secure module 108 subtracts the first value, the train fare, from the monetary value of the portable module 102 X7. The decrypted transaction count is then incremented.

A register within the secure module 108 is increased by the amount of the first value, the train fare, so that the secure

module can keep an accounting of the amount of "money" it has collected X8. The secure module 108 creates a data packet, a second data, which comprises at least the portable module's serial number, the incremented transaction count, and the reduced monetary value of the portable module 102. The second data packet is then encrypted by the secure module 108 using a private key X9.

The microprocessor based device 104 receives the encrypted second data packet, passes the encrypted second data packet to the portable module 102 X10, and opens the turn style to let the module's user onto the train platform. The portable module 102 receives the encrypted second data packet and stores it in memory X11. The portable module also increments its transaction count indicating that another transaction has occurred X12.

Thus, the above description indicates how valuable information can be transferred between a portable insecure module 102 and a secure module 108 wherein there is a conservation of value. That is, no value is gained or lost. Value that was in the portable module 102 was decreased by the same amount value was added to the secure module 108. In the example provided, the decrease and increase in value was equal to a train fare. Such an increment or decrement can also be equal to an amount provided by an ATM, credit card transaction, cash acceptor, etc.

It is also understood that the insecure portable is module 102 could be another secure module similar to the secure module in the system, but programed to act like a portable module 102.

### B. Transferring Units of Exchange Into the Portable Module 102

In this example, for simplicity, suppose the portable module does not have any monetary value and the user of the portable module wishes to "fill it up" with value. Suppose the user wishes to take cash out of an ATM machine and instead of pocketing the cash, the user wishes to put the cash value into the portable module 102.

Referring to FIG. 5, the portable module 102 contains its ID number, a transaction count and an encrypted data packet containing the portable module's ID number, transaction count and the monetary value of the portable module 102 Y1. The microprocessor based device 104, which in this example could be part of the ATM machine 112, receives the information contained in the portable module 102 when a communication is initiated between the portable module 102 and the microprocessor based device 104 Y2.

The microprocessor based device 104 passes the module's serial number, transaction count, and encrypted data packet as a first data packet to the secure module 108. The microprocessor based device also passes the amount of amount of monetary value to add to the portable module 102, as indicated by the ATM 112, to the secure module 108 Y3.

The secure module 108 decrypts the encrypted data passed to it using a public key Y4. The secure module 108 then makes a few comparisons to make sure that the data it has just received is valid and not counterfeit. The secure module 108 compares the serial number (ID number) received in the first data packet with the serial number (ID number) found in the decrypted data Y5. The secure module 108 also compares the transaction count passed the first data packet with the transaction count found in the decrypted data Y6. If the serial numbers and transaction counters match, then the secure module decides that the data received is valid and the secure module adds the monetary value, indicated by the ATM to the monetary value of the decrypted data Y7. The decrypted transaction count is incremented Y8. A register within the secure module may be decremented by the

same amount that the monetary value of the decrypted data was increased Y8.

The secure module 108 creates a second data packet, that contains the portable module's ID number, the incremented transaction counter and the increased monetary value. The second data packet is then encrypted using a private key Y10.

The microprocessor based device 104 reads the encrypted second data packet and sends it to the portable module 102 Y11. The portable module receives the encrypted second data packet and stores it in memory Y12. The portable module also advances its transaction counter Y13. The result being that the portable module now has the value of the cash withdrawn from the ATM 112. Furthermore, a record of the transaction may have been recorded and kept in the secure module, as well as by the bank that operates the ATM 112. Exemplary Firmware Definitions for Use With the Secure Module

**Object** The most primitive data structure accepted by and operated on by the secure modules firmware. A list of valid objects and their definitions is provided in the next section.

**Group** A self-contained collection of objects. An object's scope is restricted to the group of which it is a member.

**Group ID** A number preferably between 0 and 255 representing a specific group.

**Object ID** A number preferably between 0 and 255 representing a specific object within a specific group.

**Object Type** Preferably a 1-byte type specifier that describes a specific object.

**PIN** An alphanumeric Personal Identification number that is preferably eight bytes in length.

**Common PIN** The PIN that controls access to shared resources such as the audit trail. It is also used to control the host's ability to create and delete groups.

**Group PIN** The PIN that controls access to operations specific to objects within a group.

**Audit Trail** A record of transactions occurring after the secure module has been locked.

**Locked Object** An object which has been locked by executing the lock object command. Once an object is locked it is not directly readable.

**Private Object** An object which has been privatized by executing the privatize object command. Once an object is private, it is not directly readable or writable.

**Locked Group** A group which has been locked using the locked group command. After a group has been locked it will not allow object creation.

**Composite Object** A combination of several objects. The individual objects inherit the attributes of the composite object.

#### Exemplary Object Definitions

**RSA Modulus** A large integer preferably of at most 1024 bits in length. It is the product of 2 large prime numbers that are each about half the number of bits in length of the desired modulus size. The RSA modulus is used in the following equations for encrypting and decrypting a message M:

$$\text{Encryption: } C=M^e \pmod{N} \quad (1)$$

$$\text{Decryption: } M=C^d \pmod{N} \quad (2)$$

where C is the cyphertext, d and e are the RSA exponents (see below), and N is the RSA modulus.

**RSA Exponent** Both e and d (shown in equations 1 and 2 above) are RSA exponents. They are typically large numbers but are smaller than the modulus (N). RSA

exponents can be either private or public. When RSA exponents are created in the secure module, they may be declared as either. Once created an exponent may be changed from a public exponent to a private exponent. After an exponent has been made private, however, it will remain private until the transaction group 40 to which it belongs is destroyed.

**Transaction Script** A transaction script is a series of instructions to be carried out by the secure module. When invoked the secure module firmware interprets the instructions in the script and places the results in the output data object (see below). The actual script is simply a list of objects. The order in which the objects are listed specifies the operations to be performed on the objects. transaction scripts 44 preferably may be as long as 128 bytes.

**Transaction Counter** The transaction counter object is preferably 4 bytes in length and is usually initialized to zero when it is created. Every time a transaction script, which references this object, is invoked, the transaction counter increments by 1. Once a transaction counter has been locked it is read only and provides an irreversible counter.

**Money Register** The money register object is preferably 4 bytes in length and may be used to represent money or some other form of credit. Once this object has been created, it must be locked to prevent a user from tampering with its value. Once locked the value of this object can be altered only by invoking a transaction script. A typical transaction group 40 which performs monetary transactions might have one script for withdrawals from the money register and one for deposits to the money register.

**Clock Offset** This object is preferably a 4 byte number which contains the difference between the reading of the secure module's real-time clock and some convenient time (e.g., 12:00 a.m., Jan. 1, 1970). The true time can then be obtained from the secure module by adding the value of the clock offset to the real-time clock.

**SALT** A SALT object is preferably 20 bytes in length and should be initialized with random data when it is created. When a host transmits a generate random SALT command, the secure module combines the previous SALT with the secure module's random number (produced preferably by randomly occurring power-ups) to generate a new random SALT. If the SALT object has not been privatized it may subsequently be read by issuing a read object command.

**Configuration Data** This is a user defined structure with preferably a maximum length of 128 bytes. This object is typically used to store configuration information specific to its transaction group 40. For example, the configuration data object may be used to specify the format of the money register object (i.e., the type of currency it represents). Since this object has no pre-defined structure, it may never be used by a transaction object.

**Input Data** An input data object is simply an input buffer with preferably a maximum length of 128 bytes. A transaction group may have multiple input objects. The host uses input data objects to store data to be processed by transaction scripts 44.

**Output Data** The output data object is used by transaction scripts as an output buffer. This object is automatically created when the transaction group is created. It is preferably 512 bytes in length and inherits password protection from its group.

**Random Fill** When the script interpreter encounters this type of object it automatically pads the current message so that its length is 1 bit smaller than the length of the preceding

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modulus. A handle to this object is automatically created when the transaction group is created. It is a private object and may not be read using the read object command.

**Working Register** This object is used by the script interpreter as working space and may be used in a transaction script. A handle to this object is automatically created when the transaction group is created. It is a private object and may not be read using the read object command.

**ROM Data** This object is automatically created when the transaction group is created. It is a locked object and may not be altered using the write object command. This object is 8 bytes and length and its contents are identical to the 8 by ROM data of the Micro-In-A-Can™.

**Preferred Secure Module Firmware Command Set**

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Set Common PIN(01H)

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Transmit (to secure module)	01H, old PIN, new PIN, PIN option byte
Receive data	CSB (command status byte) = 0 if successful, appropriate error code otherwise
	Output length = 0
	Output Data = 0

---

Notes:

The PIN option byte may be the bitwise-or of any of the following values:

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PIN__IO_ERASE	00000001b (require PIN for Master Erase)
PIN__TO_CREATE	00000010b (require PIN for group creation).

---

Initially the secure module has a PIN (Personal Identification Number) of 0 (Null) and an option byte of 0. Once a PIN has been established it can only be changed by providing the old PIN or by a Master Erase. However, if the PIN\_\_TO\_ERASE bit is set in the option byte, the PIN can only be changed through the set common PIN command.

Possible error codes for the set common PIN command:

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ERR_BAD_COMMON_PIN	(Common PIN match failed)
ERR_BAD_PIN_LENGTH	(New PIN length > 8 bytes)
ERR_BAD_OPTION_BYTE	(Unrecognizable option byte)

---

For all commands described in this section, data received by the host will be in the form of a return packet. A return packet has the following structure:

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Command status byte	(0 if command successful, error code otherwise, 1 byte)
Output data length	(Command output length, 2 bytes)
Output data	(Command output, length specified above).

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Master Erase (02H)

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Transmit data	02H, Common PIN
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Master Erase (02H)

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Receive data	CSB = 0 if command was successful, ERR_BAD_COMMON_PIN otherwise
	Output length = 0
	Output data = 0

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Notes:

If the LSB (least significant bit) of the PIN option is clear (i.e. PIN not required for Master Erase) then a 0 is transmitted for the Common PIN value. In general this text will always assume a PIN is required. If no PIN has been established a 0 should be transmitted as the PIN. This is true of the common PIN and group PINS (see below). If the PIN was correct the firmware deletes all groups (see below) and all objects within the groups. The common PIN and common PIN option byte are both reset to zero.

After everything has been erased the secure module transmits the return packet. The CSB is as described above. The output data length and output data fields are both set to 0.

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Create Group (03H)

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Transmit data	03H, Common PIN, Group name, Group PIN
Receive data	CSB = 0 if command successful, appropriate error code otherwise
	Output length = 1 if successful, 0 otherwise
	Output data = Group ID if successful, 0 otherwise

---

Notes:

The maximum group name length is 16 bytes and the maximum PIN length is eight bytes. If the PIN\_\_TO\_CREATE bit is set in the common PIN option byte and the PIN transmitted does not match the common PIN the secure module will set the OSC to ERR\_BAD\_COMMON\_PIN.

Possible error return codes for the create group command:

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ERR_BAD_COMMON_PIN	(Incorrect common PIN)
ERR_BAD_NAME_LENGTH	(If group name length > 16 bytes)
ERR_BAD_PIN_LENGTH	(If group PIN length > 8 bytes)
ERR_MIAC_LOCKED	(The secure module has been locked)
ERR_INSUFFICIENT_RAM	(Not enough memory for new group)

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Set Group PIN (04H)

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Transmit data	04H, Group ID, old GPIN, new GPIN
Receive data	CSB = 0 if command successful, appropriate error code otherwise
	Output length = 0
	Output data = 0

---

Notes:

The Group PIN only restricts access to objects within the group specified by the group ID transmitted in the command packet.

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Possible error codes for the set group PIN command:

ERR_BAD_GROUP_PIN	(Group PIN match failed)
ERR_BAD_PIN_LENGTH	(New group PIN length > 8 bytes)

Create Object (05H)

Transmit data  
05H, Group ID, Group PIN, Object type, Object attributes, Object data  
Receive data  
CSB = 0 if command successful, appropriate error code otherwise  
Output length = 1 if successful, 0 otherwise  
Output data = object ID if successful, 0 otherwise

Notes:

If the Create Object command is successful the secure module firmware returns the object's ID within the group specified by the Group ID. If the PIN supplied by the host was incorrect or the group has been locked by the Lock Group command (described below) the secure module returns an error code in the CSB. An object creation will also fail if the object is invalid for any reason. For example, if the object being created is an RSA modulus (type 0) and it is greater than 1024 bits in length. transaction script creation will succeed if it obeys all transaction scripts rules.

Possible error return codes for the create object command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_GROUP_LOCKED	(The group has been locked)
ERR_MIAC_LOCKED	(The secure module has been locked)
ERR_INVALID_TYPE	(The object type specified is invalid)
ERR_BAD_SIZE	(The objects length was invalid)
ERR_INSUFFICIENT_RAM	(Not enough memory for new object)
Object types:	
	RSA modulus 0
	RSA exponent 1
	Money register 2
	Transaction counter 3
	Transaction script 4
	Clock offset 5
	Random SALT 6
	Configuration object 7
	Input data object 8
	Output data object 9
Object Attributes:	
	Locked 0000001b
	Privatized 0000010b

Objects may also be locked and privatized after creation by using the Lock Object and Privatize Object commands described below.

Lock Object (06H)

Transmit data  
06H, Group ID, Group PIN, Object ID  
Receive data  
CSB = 0 if command successful, appropriate

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-continued

Lock Object (06H)

error code otherwise  
Output length = 0  
Output data = 0

Notes:

If the Group ID, Group PIN and Object ID are all correct, the secure module will lock the specified object. Locking an object is an irreversible operation.

Possible error return codes for the lock object command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_GROUP_LOCKED	(The group has already been locked)
ERR_MIAC_LOCKED	(The secure module has been locked)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_OBJECT_ID	(Specified object does not exist)

Privatize Object (07H)

Transmit data  
07H, Group ID, Group PIN, Object ID  
Receive data  
CSB = 0 if successful, appropriate error code otherwise

Notes:

If the Group ID, Group PIN and Object ID were valid the object will be privatized. Privatized objects share all the properties of locked objects but are not readable. Privatized objects are only modifiable through transaction scripts. Note that locking a privatized object is legal, but has no meaning since object privatization is a stronger operation than object locking. Privatizing an object is an irreversible operation.

Possible error return codes for the privatize object command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_GROUP_LOCKED	(The group has already been locked)
ERR_MIAC_LOCKED	(The secure module has been locked)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_OBJECT_ID	(Specified object does not exist)

Make Object Destructable (08H)

Transmit data  
08H, Group ID, Group PIN, Object ID  
Receive data  
CSB = 0 if successful, appropriate error code otherwise

Notes:

If the Group ID, Group PIN and Object ID were valid the object will be made destructable. If an object is destructable it becomes unusable by a transaction script after the groups destructor becomes active. If no destructor object exists within the transaction group the destructible object attribute

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bit has no affect. Making an object destructable is an irreversible operation.

Possible error return codes for the make object destructable command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_GROUP_LOCKED	(The group has already been locked)
ERR_MIAC_LOCKED	(The secure module has been locked)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_OBJECT_ID	(Specified object does not exist)

Lock Secure module (09H)	
Transmit data	09H, Common PIN
Receive data	CSB = 0 if successful, appropriate error code otherwise
	Output length = 2 if successful, 0 otherwise
	Output data = audit trail size if successful, 0 otherwise

Notes:

If the host supplied Common PIN is correct and the secure module has not previously been locked, the command will succeed. When the secure module is locked it will not accept any new groups or objects. This implies that all groups are automatically locked. The RAM not used by the system or by groups will be used for an audit trail. There is no audit trail until the secure module has successfully been locked!

An audit trail record is six bytes long and has the following structure:

Group ID|Object ID|Date/Time stamp.

Once an audit trail has been established, a record of the form shown above will be stored in the first available size byte location every time a transaction script is executed. Note that since the secure module must be locked before the audit trail begins, neither the group ID nor any object ID is subject to change. This will always allow an application processing the audit trail to uniquely identify the transaction script that was executed. Once the audit trail has consumed all of its available memory, it will store new transaction records over the oldest transaction records.

Possible error codes for the lock secure module command:

ERR_BAD_COMMON_PIN	(Supplied common PIN was incorrect)
ERR_MIAC_LOCKED	(Secure module was already locked)

Lock Group (0AH)	
Transmit data	0AH, Group ID, Group PIN
Receive data	CSB = 0 if command successful, appropriate

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-continued

Lock Group (0AH)	
error code otherwise	Output length = 0
	Output data = 0

Notes:

If the group PIN provided is correct the secure module BIOS will not allow further object creation within the specified group. Since groups are completely self-contained entities they may be deleted by executing the Delete Group command (described below).

Possible error return codes for the lock group command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_GROUP_LOCKED	(The group has already been locked)
ERR_MIAC_LOCKED	(The secure module has been locked)
ERR_BAD_GROUP_ID	(Specified group does not exist)

Invoke Transaction Script (0BH)	
Transmit data	0BH, Group ID, Group PIN, Object ID
Receive data	CSB = 0 if command successful, appropriate error code otherwise
	Output length = 1 if successful, 0 otherwise
	Output data = estimated completion time

Notes:

The time estimate returned by the secure module is in sixteenths of a second. If an error code was returned in the CSB, the time estimate will be 0.

Possible error return codes for the execution transaction script command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_OBJECT_ID	(Script object did not exist in group)

Read Object (0CH)	
Transmit data	0CH, Group ID, Group PIN, Object ID
Receive data	CSB = 0 if command successful, appropriate error code otherwise
	Output length = object length if successful, 0 otherwise
	Output data = object data if successful, 0 otherwise

Notes:

If the Group ID, Group PIN and Object ID were correct, the secure module checks the attribute byte of the specified object. If the object has not been privatized the secure module will transmit the object data to the host. If the Group PIN was invalid or the object has been privatized the secure module will return a 0 in the output length, and data fields of the return packet.

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Possible error codes for the read object command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_OBJECT_ID	(Object did not exist in group)
ERR_OBJECT_PRIVATIZED	(Object has been privatized)

Write Object (0DH)

Transmit data	0DH, Group ID, Group PIN, Object ID, Object size, Object Data
Receive data	CSB = 0 if successful, appropriate error code otherwise
Output length = 0	Output data = 0

Notes:

If the Group ID, Group PIN and Object ID were correct, the secure module checks the attribute byte of the specified object. If the object has not been locked or privatized the secure module will clear the objects previous size and data and replace it with the new object data. Note that the object type and attribute byte are not affected.

Possible error codes for the write object command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_OBJECT_ID	(Object did not exist in group)
ERR_BAD_OBJECT_SIZE	(Illegal object size specified)
ERR_OBJECT_LOCKED	(Object has been locked)
ERR_OBJECT_PRIVATIZED	(Object has been privatized)

Read Group Name (0EH)

Transmit data	0EH, Group ID
Receive data	CSB = 0
Output Length = length of group name	Output data = group name

Notes:

The group name length is a maximum of 16 bytes. All byte values are legal in a group name.

Delete Group (0FH)

Transmit data	0FH, Group ID, Group PIN
Receive data	CSB = 0 if successful, appropriate error code otherwise
Output length = 0	Output data = 0

Notes:

If the group PIN and group ID are correct the secure module will delete the specified group. Deleting a group

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causes the automatic destruction of all objects within the group. If the secure module has been locked the Delete Group command will fail.

Possible error codes for the delete group command:

ERR_BAD_CROUP_PIN	(Incorrect group PIN)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_MIAC_LOCKED	(Secure module has been locked)

Get Command Status Info (10H)

Transmit data	10H
Receive data	CSB = 0
Output length = 6	Output data = secure module status structure (see below)

Notes:

This operation requires no PIN and never fails. The status structure is defined as follows:

Last command executed	(1 byte)
Last command status	(1 byte)
Time command received	(4 bytes)

Get Secure module Configuration Info (11H)

Transmit data	11H
Receive data	CSB = 0
Output length = 4	Output data = secure module configuration structure

Notes:

This operation requires no PIN and never fails. The configuration structure is defined as follows:

Number of groups	(1 byte)
Flag byte (see below)	(1 byte)
Audit trail size/Free RAM	(2 bytes)

The flag byte is the bitwise-or of any of the following values:

0000001b	(Secure module is locked)
0000010b	(Common PIN required for access)

Read Audit Trail Info (12H)

Transmit data	12H, Common PIN
Receive data	CSB = 0 if command successful, appropriate error code otherwise
Output length = audit trail structure size (5)	if successful, 0 otherwise



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Read Audit Trail Info (12H)	
Output data = audit trail info structure if successful, 0 otherwise	5

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Notes:

If the transmitted Common PIN is valid and the secure module has been locked, it returns audit trail configuration information as follows:

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Number of used transaction records	(2 bytes)
Number of free transaction records	(2 bytes)
A boolean specifying whether or not the audit trail rolled since previous read command	(1 byte)

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Possible error codes for the read audit trail info command:

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ERR_BAD_COMMON_PIN	(Common PIN was incorrect)
ERR_MIAC_NOT_LOCKED	(Secure module is not locked)

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Read Audit Trail (13H)	
Transmit data	13H, Common PIN
Receive data	CSB = 0 if command successful, appropriate error code otherwise
	Output length = # of new records * 6 if successful, 0 otherwise
	Output data = new audit trail records

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Notes:

If the transmitted common PIN is valid and the secure module has been locked, it will transfer all new transaction records to the host.

Possible error codes for the read audit trail command:

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ERR_BAD_COMMON_PIN	(Common PIN was incorrect).
ERR_MIAC_NOT_LOCKED	secure module is not locked

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Read Group Audit Trail (14H)	
Transmit data	14H, Group ID, Group PIN
Receive data	CSB = 0 if command successful, appropriate error code otherwise
	Output length = # of records for group * 6 if successful, 0 otherwise
	Output data = audit trail records for group

---

Notes:

This command is identical to the read audit trail command, except that only records involving the group ID specified in the transmit data are returned to the host. This allows transaction groups to record track their own activities without seeing other groups records.

Possible error codes for the read group audit trail command:

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ERR_BAD_GROUP_ID	(Group ID does not exist)
ERR_PAD_GROUP_PIN	(Common PIN was incorrect)
ERR_MIAC_NOT_LOCKED	(The secure module is not locked)

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Read Real Time Clock (15H)	
Transmit data	15H, Common PIN
Receive data	CSB = 0 if the common PIN matches and ERR_BAD_COMMON_PIN otherwise
	Output length = 4
	Output data = 4 most significant bytes of the real time clock

---

Notes:

This value is not adjusted with a clock offset. This command is normally used by a service provider to compute a clock offset during transaction group creation.

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Read Real Time Clock Adjusted (16H)	
Transmit data	16H, Group ID, Group PIN, ID of offset object
Receive data	CSB = 0 if successful, appropriate error code otherwise
	Output length = 4 if successful, 0 otherwise
	Output data = Real time clock + clock offset ID

---

Notes:

This command succeeds if the group ID and group PIN are valid, and the object ID is the ID of a clock offset. The secure module adds the clock offset to the current value of the 4 most significant bytes of the RTC and returns that value in the output data field. Note that a transaction script may be written to perform the same task and put the result in the output data object.

Possible error codes for the real time clock adjusted command:

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ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_OBJECT_TYPE	(Object ID is not a clock offset)

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Get Random Data (17H)	
Transmit data	17H, Length (L)
Receive data	CSB = 0 if successful, appropriate error code otherwise
	Output length = L if successful, 0 otherwise
	Output data = L bytes of random data if successful

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Notes:

This command provides a good source of cryptographically useful random numbers.

Possible error codes for the get random data command are:

ERR_BAD_SIZE	(Requested number of bytes > 128)
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Get Firmware Version ID (18H)

Transmit data	18H
Receive data	CSB = 0
Output length = Length of firmware version ID string	
Output data = Firmware version ID string	

Notes:

This command returns the firmware version ID as a Pascal type string (length+data).

Get Free RAM (19H)

Transmit data	19H
Receive data	CSB = 0
Output length = 2	
Output data = 2 byte value containing the amount of free RAM	

Notes:

If the secure module has been locked the output data bytes will both be 0 indicating that all memory not used by transaction groups has been reserved for the audit trail.

Change Group Name (1AH)

Transmit data	1AH, Group ID, Group PIN, New Group name
Receive data	CSB = 0 if successful or an appropriate error code otherwise
Output length = 0	
Output data = 0	

Notes:

If the group ID specified exists in the secure module and the PIN supplied is correct, the transaction group name is replaced by the new group name supplied by the host. If a group ID of 0 is supplied the PIN transmitted must be the common PIN. If it is correct, the secure module name is replaced by the new name supplied by the host.

Possible error codes for the change group name command:

ERR_BAD_GROUP_PIN	(Incorrect group PIN)
ERR_BAD_GROUP_ID	(Specified group does not exist)
ERR_BAD_NAME_LENGTH	(New group name > 16 bytes)

ERROR CODE DEFINITIONS

ERR\_BAD\_COMMAND (80H)

This error code occurs when the secure module firmware does not recognize the command just transmitted by the host.

ERR\_BAD\_COMMON\_PIN (81H)

This error code will be returned when a command requires a common PIN and the PIN supplied does not match the secure module's common PIN. Initially the common PIN is set to 0.

ERR\_BAD\_GROUP\_PIN (82H)

Transaction groups may have their own PIN, FIG. 6. If this PIN has been set (by a set group PIN command) it must be supplied to access any of the objects within the group. If the Group PIN supplied does not match the actual group PIN, the secure module will return the ERR\_BAD\_GROUP\_PIN error code.

ERR\_BAD\_PIN\_LENGTH (83H)

There are 2 commands which can change PIN values. The set group PIN and the set common PIN commands. Both of these require the new PIN as well as the old PIN. The ERR\_BAD\_PIN\_LENGTH error code will be returned if the old PIN supplied was correct, but the new PIN was greater than 8 characters in length.

ERR\_BAD\_OPTION\_BYTE (84H)

The option byte only applies to the common PIN. When the set common PIN command is executed the last byte the host supplies is the option byte (described in command section). If this byte is unrecognizable to the secure module, it will return the ERR\_BAD\_OPTION\_BYTE error code.

ERR\_BAD\_NAME\_LENGTH (85H)

When the create transaction group command is executed, one of the data structures supplied by the host is the group's name. The group name may not exceed 16 characters in length. If the name supplied is longer than 16 characters, the ERR\_BAD\_NAME\_LENGTH error code is returned.

ERR\_INSUFFICIENT\_RAM (86H)

The create transaction group and create object commands return this error code when there is not enough heap available in the secure module.

ERR\_MIAC\_LOCKED (87H)

When the secure module has been locked, no groups or objects can be created or destroyed. Any attempts to create or delete objects will generate an ERR\_MIAC\_LOCKED error code.

ERR\_MIAC\_NOT\_LOCKED (88H)

If the secure module has not been locked there is no audit trail. If one of the audit trail commands is executed this error code will be returned.

ERR\_GROUP\_LOCKED (89H)

Once a transaction group has been locked object creation within that group is not possible. Also the objects attributes and types are frozen. Any attempt to create objects or modify their attribute or type bytes will generate an ERR\_GROUP\_LOCKED error code.

ERR\_BAD\_OBJECT\_TYPE (8AH)

When the host sends a create object command to the secure module, one of the parameters it supplies is an object

type (see command section). If the object type is not recognized by the firmware it will return an ERR\_BAD\_OBJECT\_TYPE error code.

#### ERR\_BAD\_OBJECT\_ATTR (8BH)

When the host sends a create object command to the secure module, one of the parameters it supplies is an object attribute byte (see command section). If the object attribute byte is not recognized by the firmware it will return an ERR\_BAD\_OBJECT\_ATTR error code.

#### ERR\_BAD\_SIZE (8CH)

An ERR\_BAD\_SIZE error code is normally generated when creating or writing an object. It will only occur when the object data supplied by the host has an invalid length.

#### ERR\_BAD\_GROUP\_ID (8DH)

All commands that operate at the transaction group level require the group ID to be supplied in the command packet. If the group ID specified does not exist in the secure module it will generate an ERR\_BAD\_GROUP\_ID error code.

#### ERR\_BAD\_OBJECT\_ID (8EH)

All commands that operate at the object level require the object ID to be supplied in the command packet. If the object ID specified does not exist within the specific transaction group (also specified in the command packet) the secure module will generate an ERR\_BAD\_OBJECT\_ID error code.

#### ERR\_INSUFFICIENT\_FUNDS (8FH)

If a script object that executes financial transactions is invoked and the value of the money register is less than the withdrawal amount requested an ERR\_INSUFFICIENT\_FUNDS error code will be returned.

#### ERR\_OBJECT\_LOCKED (90H)

Locked objects are read only. If a write object command is attempted and it specifies the object ID of a locked object the secure module will return an ERR\_OBJECT\_LOCKED error code.

#### ERR\_OBJECT\_PRIVATE (91H)

Private objects are not directly readable or writable. If a read object command or a write object command is attempted, and it specifies the object ID of a private object, the secure module will return an ERR\_OBJECT\_PRIVATE error code.

#### ERR\_OBJECT\_DESTROYED (92H)

If an object is destructible and the transaction group's destructor is active the object may not be used by a script. If a script is invoked which uses an object which has been destructed, an ERR\_OBJECT\_DESTROYED error code will be returned by the secure module.

The exemplary embodiment of the present invention is preferably placed within a durable stainless steel, token-like can. It is understood that an exemplary secure module can be placed in virtually any articulatable item. Examples of

articulatable items include credit cards, rings, watches, wallets, purses, necklaces, jewelry, ID badges, pens, clipboards, etc.

The secure module 108 preferably is a single chip "trusted computer". By the word "trusted" it is meant that the computer is extremely secure from tampering by unwarranted means. The secure module incorporates a numeric coprocessor optimized for math intensive encryption. The BIOS is preferably immune to alteration and specifically designed for very secure transactions.

Each secure module can have a random "seed" generator with the ability to create a private/public key set. The private key never leaves the secure module and is only known by the secure module. Furthermore, discovery of the private key is prevented by active self-destruction upon wrongful entry into the secure module. The secure module can be bound to the user by a personal identification number (PIN).

When transactions are performed by the secure module 108 certificates of authentication are created by either or both the secure module and a system the secure module communicates with. The certificate can contain a variety of information. In particular, the certificate may contain:

- 1) who is the secure module user via a unique registration number and a certified public key.
- 2) when the transaction took place via a true-time stamping of the transaction.
- 3) where the transaction took place via a registered secure module interface site identification.
- 4) security information via uniquely serialized transactions and digital sign on message digests.
- 5) secure module status indicated as valid, lost, or expired.

Although a preferred embodiment of the method and apparatus of the present invention has been illustrated in the accompanying Drawings and described in the foregoing Detailed Description, it will be understood that the invention is not limited to the embodiment disclosed, but is capable of numerous rearrangements, modifications and substitutions without departing from the spirit of the invention as set forth and defined by the following claims.

What is claimed is:

1. A method for electronically transferring units of exchange between a first module and a second module, comprising the steps of:

- a. initiating communication between said first module and an electronic device;
- b. passing a first value datum from said first module to said electronic device;
- c. passing said first value datum from said electronic device to said second module;
- d. performing a mathematical calculation on said first value datum thereby creating a second value datum;
- e. passing said second value datum from said second module to said electronic device;
- f. passing said second value datum from said electronic device to said first module;
- g. storing said second value datum in said first module; and
- h. discontinuing communication between said first module and said electronic device.

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2. The method of claim 1, wherein said first value datum represents a monetary equivalent.

3. The method of claim 1, wherein said first value datum is encrypted.

4. The method of claim 1, wherein said second value datum is encrypted.

5. The method of claim 3, wherein the step of performing a mathematical calculation comprises the steps of:

- m. decrypting said first value datum with a public key thereby creating a decrypted value;

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n. performing at least one of an addition function and a subtraction function on said decrypted value thereby creating a value result; and

o. encrypting said value result with a private key thereby creating said second value datum.

6. The method of claim 1, wherein the step (b) of passing is performed over at least a single conductive contact.

\* \* \* \* \*

UNITED STATES PATENT AND TRADEMARK OFFICE  
**CERTIFICATE OF CORRECTION**

PATENT NO. : 5,949,880  
DATED : Sep. 7, 1999  
INVENTOR(S) : Curry et al.

Page 1 of 2

It is certified that error appears in the above-identified patent and that said Letters Patent is hereby corrected as shown below:

Column 2, line 57	Replace "electromagnetic" With --electro-magnetic--
Column 5, line 15	Before "information" Remove --is--
Column 8, line 26	Before "module" Remove --is--
Column 12, line 47	Replace "ERR_BAD_PIN_LENGTH" With --ERR_BAD_PIN_LENGTH--
Column 17, line 34	Replace "ERR_BAD_OBJECT_ID" With --ERR_BAD_OBJECT_ID--

UNITED STATES PATENT AND TRADEMARK OFFICE  
**CERTIFICATE OF CORRECTION**

PATENT NO. : 5,949,880  
DATED : Sept 7, 1999  
INVENTOR(S) : Curry et al

Page 2 of 2

It is certified that error appears in the above-identified patent and that said Letters Patent is hereby corrected as shown below:

Column 20, line 6	Replace "ERR MIAC NOT LOCKED" With --ERR_MIAC_NOT_LOCKED--
Column 20, line 48	Replace "ERR_BAD OBJECT TYPE" With --ERR_BAD_OBJECT_TYPE--
Column 21, line 58	Replace "ERR_BAD NAME LENGTH" With --ERR_BAD_NAME_LENGTH--

Signed and Sealed this  
Twenty-fifth Day of April, 2000

Attest:



Q. TODD DICKINSON

Attesting Officer

Director of Patents and Trademarks